

CS 302, Theory of Computation  
Even Semester, 2004-2005  
Home Assignment # 1 (Alternative)  
Due Date: 15/02/2005

07/02/2005

1. Assume that the set of atoms (and hence formulas) in propositional logic is countable, i.e. there is an enumeration  $A_1, A_2, \dots$  of all formulas. A set  $\Phi$  of formulas is called *consistent* if every finite subset of  $\Phi$  is satisfiable.  $\Phi$  is *maximally consistent* if it is consistent and if no proper superset of  $\Phi$  is consistent. Define by recursion the following sequence of sets of formulas, where  $\Phi$  is a consistent set:

$$\Delta_0 = \Phi, \quad \Delta_{n+1} = \begin{cases} \Delta_n \cup \{A_{n+1}\} & \text{if this set is consistent} \\ \Delta_n \cup \{\neg A_{n+1}\} & \text{otherwise} \end{cases}$$

Prove the following results without using the compactness theorem.

- (a) Each  $\Delta_n$  is consistent. (4 points)
- (b)  $\Delta = \bigcup_n \Delta_n$  is consistent. (6 points)
- (c) (Lindenbaum's Theorem) Show that every consistent set  $\Phi$  can be extended to a maximally consistent set. (10 points)