

Lecture 8: Drawbacks in Elimination Method

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LECTURE 8

08-AUGUST-2012

Last class we discussed on

- Banded Matrix and Sparse Matrix
- Thomas Algorithm

Drawbacks in Elimination Methods

Main drawbacks are:

- Cumulate Round-off errors
- Not suitable for ill-conditioned systems.

Round-off Errors

- * Computers can store numbers with finite precision only (say single precision with 7 significant digits, double precision with 14 significant digits, etc).
- * If you have infinite precision → then round-off errors will not appear.
- * Due to approximating infinite precision numbers using finite-precision, we get round-off errors.

Consider the linear algebraic system

$$\begin{aligned} 0.0003 x_1 + 3 x_2 &= 1.0002 \\ x_1 + x_2 &= 1 \end{aligned}$$

(2)

Using traditional Gauss elimination,

$$\left[\begin{array}{cc|c} 0.0003 & 3 & 1.0002 \\ 1 & 1 & 1 \end{array} \right] R_2 - \left(\frac{1}{0.0003} \right) R_1 \Rightarrow \left[\begin{array}{cc|c} 0.0003 & 3 & 1.0002 \\ 0 & -9999 & -3333 \end{array} \right]$$

The actual solutions should be

$$x_2 = 0.333333 \dots \dots \quad \text{and} \quad x_1 = 0.666666 \dots \dots$$

However in computer you have to restrict to finite precision.

You know $x_2 = \frac{1 - \frac{1.0002}{0.0003}}{9999}$ and $x_1 = \frac{1.0002 - 3x_2}{0.0003}$

Let us compare the results if you are using varying significant figures from 4 to 7.

Significant Fig	x_2	x_1
4	0.3333	1.0000
5	0.33333	0.7000
6	0.333333	0.6700
7	0.3333333	0.6670

∴ Based on your precision, you can see the variation in solution. This can be overcome by partial pivoting.

$$\left[\begin{array}{cc|c} 1 & 1 & 1.0002 \\ 0.0003 & 3 & 1.0002 \end{array} \right] R_2 = R_2 - 0.0003 R_1 \Rightarrow \left[\begin{array}{cc|c} 1 & 1 & 1.0002 \\ 0 & 2.9997 & 0.9999 \end{array} \right]$$

Sign Fig	x_2	x_1
4	0.3333	0.6667
5	0.33333	0.66667
6	0.333333	0.666667
7	0.3333333	0.6666667

$$(x_1 = 1 - x_2)$$

(3)

System Condition

- Usually well-posed non-singular systems have exact solutions.
- As in computers you can store only finite precision numbers, chances of round-off errors are more.
- Systems that are sensitive to small changes in numbers due to round-off errors may give erroneous solutions by elimination methods. They are not suitable.
- A well-conditioned problem is one in which a small change in any of the elements causes only small change in the solution of the problem.
- An ill-conditioned problem is the one where small changes in numerical values causes large change in solution. (Usually ill-conditioned systems are sensitive to round-off errors).

See the example as referred from Hoffman (2001) below:

$$\begin{aligned} x_1 + x_2 &= 2 \\ x_1 + 1.0001 x_2 &= 2.0001 \end{aligned}$$

$$\text{i.e. } \left[\begin{array}{cc|c} 1 & 1 & 2 \\ 1 & 1.0001 & 2.0001 \end{array} \right] R_2 = R_2 - \left(\frac{1}{1}\right)R_1 \Rightarrow \left[\begin{array}{cc|c} 1 & 1 & 2 \\ 0 & 0.0001 & 0.0001 \end{array} \right]$$

$$\therefore x_2 = 1 \text{ and } x_1 = 1$$

If we slightly modify the system as:

$$\begin{aligned} x_1 + x_2 &= 2 \\ x_1 + 0.9999 x_2 &= 2.0001 \end{aligned}$$

$$\text{i.e. } \left[\begin{array}{cc|c} 1 & 1 & 2 \\ 1 & 0.9999 & 2.0001 \end{array} \right] R_2 = R_2 - \left(\frac{1}{1}\right)R_1 \Rightarrow \left[\begin{array}{cc|c} 1 & 1 & 2 \\ 0 & -0.0001 & 0.0001 \end{array} \right]$$

$$\therefore x_2 = -1, \text{ and } x_1 = 3$$

(2)

There is quite a significant change in the solution of the given system, ~~and~~ even though only slight variation is made in the system. This is because the system is ill-conditioned.

- * With infinite precision, it may be possible to overcome ill-condition. However, in computational solutions you need to take care of avoiding ill-conditioned systems.
- * Ill-conditioning can be determined by Condition Number.

Norms and Condition Number

Norms: \rightarrow It is the measure of magnitude of a matrix or vector like $[A]$, $\{x\}$, $\{b\}$, etc.

(I am avoiding $[]$, $\{ \}$ inside the norm symbols)

• $\|A\|$, $\|x\|$, $\|b\|$ - etc.

Please note that

$$\|A\| > 0 \quad \rightarrow (i)$$

$$\text{If } \|A\| = 0 \quad \text{then } [A] = 0 \quad \rightarrow (ii)$$

it is due to

A scalar quantity k if it is multiplied to matrix A we get $[kA]$

$$\|kA\| = |k| \|A\| \quad \rightarrow (iii)$$

\hookrightarrow Modulus
or
Absolute

(5)

$[A], [B]$ exist

$$\|A + B\| \leq \|A\| + \|B\|$$

$$\|AB\| \leq \|A\| \|B\|$$

Norm of a scalar is its absolute value $\|k\| = |k|$

Norm of vector $\{x\} = \begin{Bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{Bmatrix} = x_i$

$$\|x\|_1 = \sum |x_i| \quad (\text{Sum of the magnitudes})$$

$$\|x\|_2 = \|x\|_e = (\sum x_i^2)^{1/2} \quad (\text{Euclidean Norm})$$

$$\|x\|_\infty = \max_{1 \leq i \leq n} |x_i| \quad \text{Maximum magnitude norm.}$$

For an $n \times n$ matrix $[A]$

$$\|A\|_1 = \max_{1 \leq j \leq n} \sum_{i=1}^n |a_{ij}| \quad \rightarrow \text{Maximum of the column sums.}$$

$$\|A\|_\infty = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}| \quad \rightarrow \text{Max. of row sums.}$$

$$\|A\|_e \rightarrow \text{Euclidean Norm} = \left(\sum_{i=1}^n \sum_{j=1}^n a_{ij}^2 \right)^{1/2}$$

Condition Number

\rightarrow It is the measure of sensitivity of the system to small changes in any of its elements.

Consider the linear system

$$[A]\{x\} = \{b\}$$

$$\therefore \|b\| \leq \|A\| \|x\| \quad (\text{Isn't it?})$$

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Let us change the R.H.S vector $\{b\}$ by a small amount $\{b\} + \{Sb\}$.

This will cause change in solution $\{x\} + \{Sx\}$

$$\therefore [A] (\{x\} + \{Sx\}) = \{b\} + \{Sb\}$$

$$\text{or } [A] \{x\} + [A] \{Sx\} = \{b\} + \{Sb\}$$

$$\text{or } [A] \{Sx\} = \{Sb\}$$

$$\text{or } \{Sx\} = [A]^{-1} \{Sb\}$$

$$\|Sx\| \leq \|A^{-1}\| \|Sb\|$$

or Multiplying on L.H.S by $\|b\|$ and RHS by $\|A\| \|x\|$

$$\therefore \|b\| \|Sx\| \leq \|A\| \|x\| \|A^{-1}\| \|Sb\|$$

$$\text{i.e. } \frac{\|Sx\|}{\|x\|} \leq \|A\| \|A^{-1}\| \frac{\|Sb\|}{\|b\|}$$

$$\text{i.e. } \frac{\|Sx\|}{\|x\|} \leq C(A) \frac{\|Sb\|}{\|b\|}$$

where $C(A) = \|A\| \|A^{-1}\|$ is called the condition number.

Smaller values of $C(A) \rightarrow$ small sensitivity of system (well conditioned)

Larger values of $C(A) \rightarrow$ large sensitivity (ill-conditioned).

⑦

Q. Check the following matrix $[A] = \begin{bmatrix} 1 & 1 \\ 1 & 1.0001 \end{bmatrix}$

~~See its~~ Find its conditioning.

$$[A] = \begin{bmatrix} 1 & 1 \\ 1 & 1.0001 \end{bmatrix}$$

$$\therefore [A]^{-1} = 10000 \begin{bmatrix} 1.0001 & -1 \\ -1 & 1 \end{bmatrix} = \begin{bmatrix} 10001 & -10000 \\ -10000 & 10000 \end{bmatrix}$$

$$\|A\|_e = 2.00005, \quad \|A^{-1}\| = 20000.50002$$

$$\therefore C(A) = \|A\| \|A^{-1}\| = \underline{\underline{40002}} \gg 1$$

All. conditional.

Quiz-2

- 1) What is meant by bandwidth of a multi-diagonal sparse matrix. What is the bandwidth of a tri-diagonal matrix?
- 2) Which is the element that gets modified during elimination of a tri-diagonal matrix?
- 3) What is condition number?