

Section 7.3 Formal Proofs in Predicate Calculus

All proof rules for propositional calculus extend to predicate calculus.

Example. ...

$k.$	$\forall x p(x)$	P
$k+1.$	$\forall x p(x) \rightarrow \exists x p(x)$	P
$k+2.$	$\exists x p(x)$	1, 2, MP
...		

But we need additional proof rules to reason with most quantified wffs. For example, suppose we want to prove that the following wff is valid.

$$\exists x \forall y p(x, y) \rightarrow \forall y \exists x p(x, y).$$

We might start with

$$\text{Proof: } 1. \exists x \forall y p(x, y) \quad P$$

But what do we do for the next line of the proof? We're stuck if we want to use proof rules. We need more proof rules.

Free to Replace

For a wff $W(x)$ and a term t we say t is *free to replace* x in $W(x)$ if $W(t)$ has the same bound occurrences of variables as $W(x)$.

Example. Let $W(x) = \exists y p(x, y)$. Then

$W(y) = \exists y p(y, y)$, so y is not free to replace x in $W(x)$.

$W(f(x)) = \exists y p(f(x), y)$, so $f(x)$ is free to replace x in $W(x)$.

$W(c) = \exists y p(c, y)$, so c is free to replace x in $W(x)$.

$W(x) = \exists y p(x, y)$, so x is free to replace x in $W(x)$.

Universal Instantiation (UI)

$\frac{\forall x W(x)}{W(t)}$ if t is free to replace x in $W(x)$.

Special cases that satisfy the restriction:

$$\frac{\forall x W(x)}{W(x)} \quad \frac{\forall x W(x)}{W(c)}$$

Existential Generalization (EG)

$\frac{W(t)}{\exists x W(x)}$ if t is free to replace x in $W(x)$.

Special cases that satisfy the restriction:

$$\frac{W(x)}{\exists x W(x)} \quad \frac{W(c)}{\exists x W(x)}$$

Existential Instantiation (EI)

If $\exists x W(x)$ occurs on some line of a proof, then $W(c)$ may be placed on any subsequent line of the proof subject to the following restrictions:

Choose c to be a new constant in the proof and such that c does not occur in the statement to be proven.

Universal Generalization (UG)

If $W(x)$ occurs on some line of a proof, then $\forall x W(x)$ may be placed on any subsequent line of the proof subject to the following restrictions:

Among the wffs used to obtain $W(x)$, x is not free in any premise and x is not free in any wff obtained by EI.

Restrictions on quantifier inference rules are necessary

Example. $\forall x \exists y p(x, y) \rightarrow \exists y \forall x p(x, y)$ is invalid. Here is an *attempted* proof.

1. $\forall x \exists y p(x, y)$ P
 2. $\exists y p(x, y)$ 1, UI
 3. $p(x, c)$ 2, EI
 4. $\forall x p(x, c)$ 3, UG (NO, x on line 3 is free in wff obtained by EI)
 5. $\exists y \forall x p(x, y)$ 4, EG
- NOT QED 1–5, CP

Example. $\exists x p(x) \rightarrow \forall x p(x)$ is invalid. Here is an *attempted* proof.

1. $\exists x p(x)$ P
 2. $p(x)$ 1, EI (NO, x is not a new constant in the proof)
 3. $\forall x p(x)$ 2, UG (NO, x on line 2 is free in wff obtained by EI)
- NOT QED 1–3, CP

Example. $\exists x p(x) \wedge \exists x q(x) \rightarrow \exists x (p(x) \wedge q(x))$ is invalid. Here is an *attempted* proof.

1. $\exists x p(x)$ P
 2. $\exists x q(x)$ P
 3. $p(c)$ 1, EI
 4. $q(c)$ 2, EI (NO, c is not a new constant in the proof)
 5. $p(c) \wedge q(c)$ 3, 4, Conj
 6. $\exists x (p(x) \wedge q(x))$ 5, EG
- NOT QED 1–6, CP

Example. $p(x) \rightarrow \forall x p(x)$ is invalid. Here is an *attempted* proof.

1. $p(x)$ P
2. $\forall x p(x)$ 1, UG (NO, x is free in a premise)
NOT QED 1, 2, CP

Example. $\forall x \exists y p(x, y) \rightarrow \exists y p(y, y)$ is invalid. Here is an *attempted* proof.

1. $\forall x \exists y p(x, y)$ P
2. $\exists y p(y, y)$ 1, UI (NO, y is not free to replace x in $\exists y p(x, y)$)
NOT QED 1, 2, CP

Example. $\forall x p(x, f(x)) \rightarrow \exists x p(x, x)$ is invalid. Here is an *attempted* proof.

1. $\forall x p(x, f(x))$ P
2. $p(x, f(x))$ 1, UI
3. $\exists x p(x, x)$ 2, EG (NO, $p(x, f(x)) \neq p(x, x)(x/t)$ for any term t)
NOT QED 1–3, CP

Example. $\forall x p(x, f(x)) \rightarrow \exists y \forall x p(x, y)$ is invalid. Here is an *attempted* proof.

1. $\forall x p(x, f(x))$ P
2. $\exists y \forall x p(x, y)$ 1, EG (NO, $f(x)$ is not free to replace y in $\forall x p(x, y)$)
NOT QED 1, 2, CP

Example. $\exists x p(x) \rightarrow p(c)$ is invalid. Here is an *attempted* proof.

1. $\exists x p(x)$ P
2. $p(c)$ 1, EI (NO, c occurs in statement to be proved)
NOT QED 1, 2, CP

Now Some Valid Wffs

Example. $\forall x \forall y p(x, y) \rightarrow \forall y p(y, y)$ is valid. Here is an *attempted* proof.

1. $\forall x \forall y p(x, y)$ P
2. $\forall y p(y, y)$ 1, UI (NO, y is not free to replace x in $\forall y p(x, y)$)
NOT QED 1, 2, CP

But here is a *correct* proof.

1. $\forall x \forall y p(x, y)$ P
2. $\forall y p(x, y)$ 1, UI
3. $p(x, x)$ 2, UI
4. $\forall x p(x, x)$ 3, UG
5. $p(y, y)$ 4, UI
6. $\forall y p(y, y)$ 5, UG
QED 1–6, CP.

Quiz. Find a proof of the statement that uses IP.

Example/Quiz. $\forall x (A(x) \rightarrow B(x)) \rightarrow (\forall x A(x) \rightarrow \forall x B(x))$ is valid. Find a proof.

1. $\forall x (A(x) \rightarrow B(x))$ P
2. $\forall x A(x)$ P [for $\forall x A(x) \rightarrow \forall x B(x)$]
3. $A(x)$ 2, UI
4. $A(x) \rightarrow B(x)$ 1, UI
5. $B(x)$ 3, 4, MP
6. $\forall x B(x)$ 5, UG
7. $\forall x A(x) \rightarrow \forall x B(x)$ 2–6, CP
QED 1, 7, CP.

Example/Quiz. Prove that the following wff is valid using IP.

$$\forall x \neg p(x, x) \wedge \forall x \forall y \forall z (p(x, y) \wedge p(y, z) \rightarrow p(x, z)) \rightarrow \forall x \forall y \neg (p(x, y) \wedge p(y, x)).$$

- | | |
|---|---|
| 1. $\forall x \neg p(x, x)$ | P |
| 2. $\forall x \forall y \forall z (p(x, y) \wedge p(y, z) \rightarrow p(x, z))$ | P |
| 3. $\exists x \exists y (p(x, y) \wedge p(y, x))$ | P [for $\forall x \forall y \neg (p(x, y) \wedge p(y, x))$], T |
| 4. $p(a, b) \wedge p(b, a)$ | 3, EI, EI |
| 5. $p(a, b) \wedge p(b, a) \rightarrow p(a, a)$ | 2, UI, UI, UI |
| 6. $p(a, a)$ | 4, 5, MP |
| 7. $\neg p(a, a)$ | 1, UI |
| 8. False | 6, 7, Contr |
| 9. $\forall x \forall y \neg (p(x, y) \wedge p(y, x))$ | 3–8, IP |
| QED | 1, 2, 9, CP. |

Quiz. Find a CP proof of the statement.

- | | |
|---|---------------|
| 1. $\forall x \neg p(x, x)$ | P |
| 2. $\forall x \forall y \forall z (p(x, y) \wedge p(y, z) \rightarrow p(x, z))$ | P |
| 3. $\neg p(x, x)$ | 1, UI |
| 4. $p(x, y) \wedge p(y, x) \rightarrow p(x, x)$ | 2, UI, UI, UI |
| 5. $\neg (p(x, y) \wedge p(y, x))$ | 3, 4, MT |
| 6. $\forall x \forall y \neg (p(x, y) \wedge p(y, x))$ | 5, UG, UG |
| QED | 1–6, CP. |

Example/Quiz. Use IP to prove that the $\forall x \exists y (p(x) \rightarrow p(y))$ is valid.

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|--|--------------|
| 1. $\exists x \forall y (p(x) \wedge \neg p(y))$ | P [for IP] |
| 2. $\forall y (p(c) \wedge \neg p(y))$ | 1, EI |
| 3. $p(c) \wedge \neg p(c)$ | 2, UI |
| 4. $p(c)$ | 3, Simp |
| 5. $\neg p(c)$ | 4, Simp |
| 6. False | 4, 5, Contr |
| QED | 1–6, IP. |

Group Quiz. Divide the class into six subgroups and assign each group one of the following six wffs to prove using CP (no IP and no T's). Assume that x does not occur free in C .

1. $\forall x (A(x) \rightarrow C) \rightarrow (\exists x A(x) \rightarrow C)$.
2. $(\exists x A(x) \rightarrow C) \rightarrow \forall x (A(x) \rightarrow C)$.
3. $(C \rightarrow \forall x A(x)) \rightarrow \forall x (C \rightarrow A(x))$.
4. $(C \rightarrow \exists x A(x)) \rightarrow \exists x (C \rightarrow A(x))$.
5. $\exists x (C \rightarrow A(x)) \rightarrow (C \rightarrow \exists x A(x))$.
6. $\exists x (A(x) \rightarrow C) \rightarrow (\forall x A(x) \rightarrow C)$.

Solutions.

1. $\forall x (A(x) \rightarrow C) \rightarrow (\exists x A(x) \rightarrow C)$.

- | | | |
|----|----------------------------------|---|
| 1. | $\forall x (A(x) \rightarrow C)$ | P |
| 2. | $\exists x A(x)$ | P [for $\exists x A(x) \rightarrow C$] |
| 3. | $A(d)$ | 2, EI |
| 4. | $A(d) \rightarrow C$ | 1, UI |
| 5. | C | 3, 4, MP |
| 6. | $\exists x A(x) \rightarrow C$ | 2–5, CP |
| | QED | 1, 6, CP. |

2. $(\exists x A(x) \rightarrow C) \rightarrow \forall x (A(x) \rightarrow C)$.

- | | | |
|----|----------------------------------|---------------------------------|
| 1. | $\exists x A(x) \rightarrow C$ | P |
| 2. | $A(x)$ | P [for $A(x) \rightarrow C$] |
| 3. | $\exists x A(x)$ | 2, EG |
| 4. | C | 1, 3, MP |
| 5. | $A(x) \rightarrow C$ | 2–4, CP |
| 6. | $\forall x (A(x) \rightarrow C)$ | 5, UG |
| | QED | 1, 5–6, CP. |

3. $(C \rightarrow \forall x A(x)) \rightarrow \forall x (C \rightarrow A(x))$.

1.	$C \rightarrow \forall x A(x)$	P
2.	C	P [for $C \rightarrow A(x)$]
3.	$\forall x A(x)$	1, 2, MP
4.	$A(x)$	3, UI
5.	$C \rightarrow A(x)$	2–4, CP
6.	$\forall x (C \rightarrow A(x))$	5, UG
	QED	1, 5–6, CP.

4. $(C \rightarrow \exists x A(x)) \rightarrow \exists x (C \rightarrow A(x))$.

1.	$C \rightarrow \exists x A(x)$	P
2.	C	P [for $C \rightarrow A(?)$]
3.	$\exists x A(x)$	1, 2, MP
4.	$A(d)$	3, EI
5.	$C \rightarrow A(d)$	2–4, CP
6.	$\exists x (C \rightarrow A(x))$	5, EG
	QED	1, 5–6, CP.

5. $\exists x (C \rightarrow A(x)) \rightarrow (C \rightarrow \exists x A(x))$.

1.	$\exists x (C \rightarrow A(x))$	P
2.	C	P [for $C \rightarrow \exists x A(x)$]
3.	$C \rightarrow A(d)$	1, EI
4.	$A(d)$	2, 3, MP
5.	$\exists x A(x)$	4, EG
6.	$C \rightarrow \exists x A(x)$	2–5, CP
	QED	1, 6, CP.

6. $\exists x (A(x) \rightarrow C) \rightarrow (\forall x A(x) \rightarrow C)$.

1.	$\exists x (A(x) \rightarrow C)$	P
2.	$\forall x A(x)$	P [for $\forall x A(x) \rightarrow C$]
3.	$A(d) \rightarrow C$	1, EI
4.	$A(d)$	2, UI
5.	C	3, 4, MP
6.	$\forall x A(x) \rightarrow C$	2–5, CP
	QED	1, 6, CP.