R.K. Bhattacharjya/CE/IITG

Introduction To Genetic Algorithms



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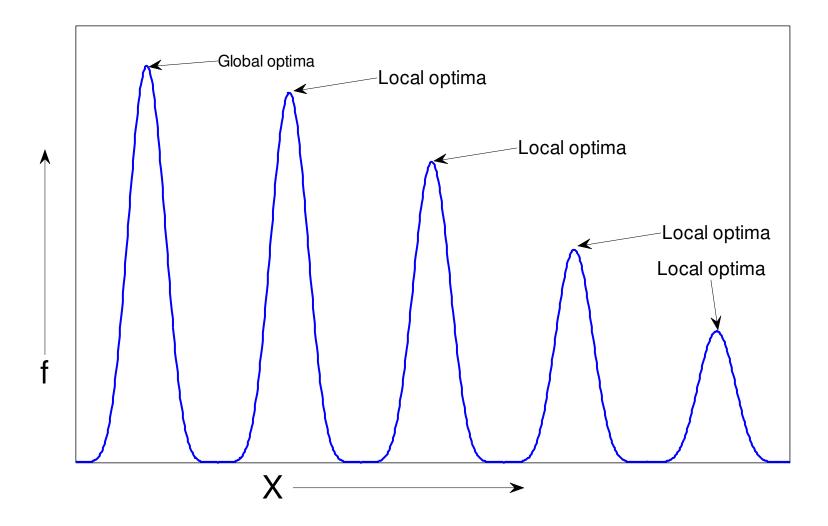
References

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- John H. Holland 'Genetic Algorithms', Scientific American Journal, July 1992.
- Kalyanmoy Deb, 'An Introduction To Genetic Algorithms', Sadhana, Vol. 24 Parts 4 And 5.

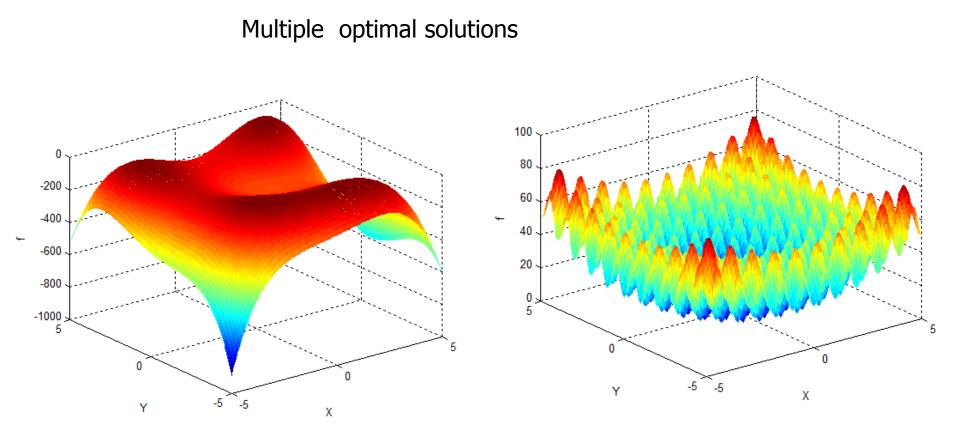
Introduction to optimization

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Introduction to optimization





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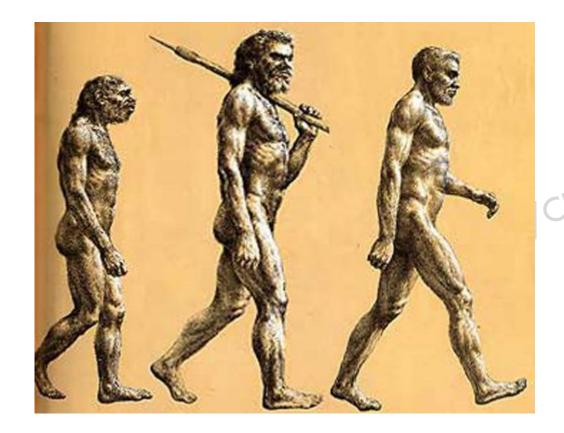
Genetic Algorithms

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Genetic Algorithms are the heuristic search and optimization techniques that mimic the process of natural evolution.

Principle Of Natural Selection

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"Select The Best, Discard The Rest"

An Example....

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Giraffes have long necks

- Giraffes with slightly longer necks could feed on leaves of higher branches when all lower ones had been eaten off.
- They had a better chance of survival.
- Favorable characteristic propagated through generations of giraffes.



Now, evolved species has long necks.

An Example....

This longer necks may have due to the effect of mutation initially. However as it was favorable, this was propagated over the generations.

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Evolution of species

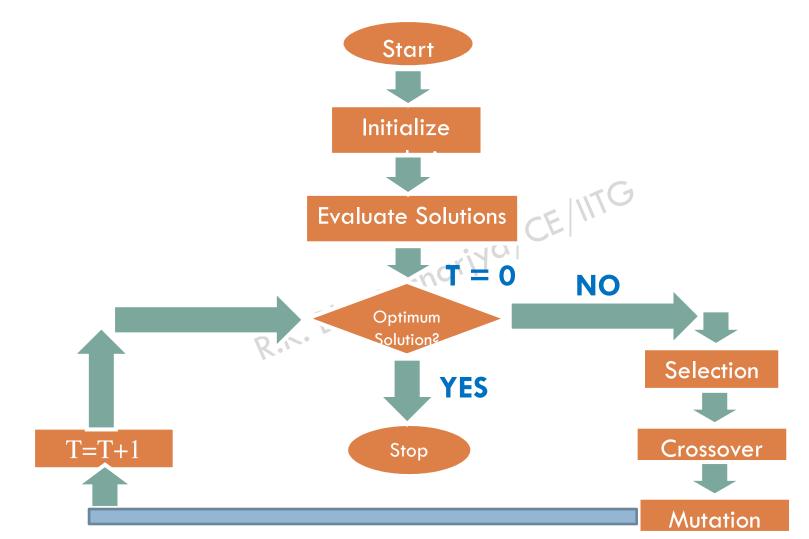
R.K. Bhattacharjya/CE/IITG Initial Population of animals INTG Struggle For Existence Millions Of Years Survival Of the Fittest 011 Surviving Individuals Reproduce, Propagate Favorable Characteristics **Evolved Species** 7 November 2013



Thus genetic algorithms implement the optimization strategies by simulating evolution of species through natural selection

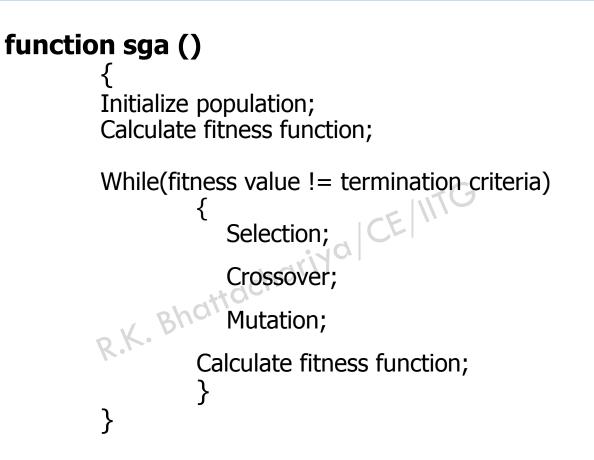
Simple Genetic Algorithms

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Simple Genetic Algorithm

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GA Operators and Parameters

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Selection

Crossover

Mutation

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Selection

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The process that determines which solutions are to be preserved and allowed to reproduce and which ones deserve to die out.

The primary objective of the selection operator is to emphasize the good solutions and eliminate the bad solutions in a population while keeping the population size constant.

"Selects the best, discards the rest"

Functions of Selection operator

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Identify the good solutions in a population

Make multiple copies of the good solutions

Eliminate bad solutions from the population so that multiple copies of good solutions can be placed in the population

Now how to identify the good solutions?

Fitness function

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A fitness value can be assigned to evaluate the solutions

A fitness function value quantifies the optimality of a solution. The value is used to rank a particular solution against all the other solutions

A fitness value is assigned to each solution depending on how close it is actually to the optimal solution of the problem

Assigning a fitness value

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Minimize $f(d, h) = c((\pi d^2/2) + \pi dh)$, Subject to $g_1(d, h) \equiv (\pi d^2 h/4) \ge 300$, Variable bounds $d_{\min} \le d \le d_{\max}$, $h_{\min} \le h \le h_{\max}$. Considering c = 0.0654 $F(s) = 0.0654(\pi(8)^2/2 + \pi(8)(10))$, = 23, d

Selection operator

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There are different techniques to implement selection in Genetic Algorithms.

□ They are:

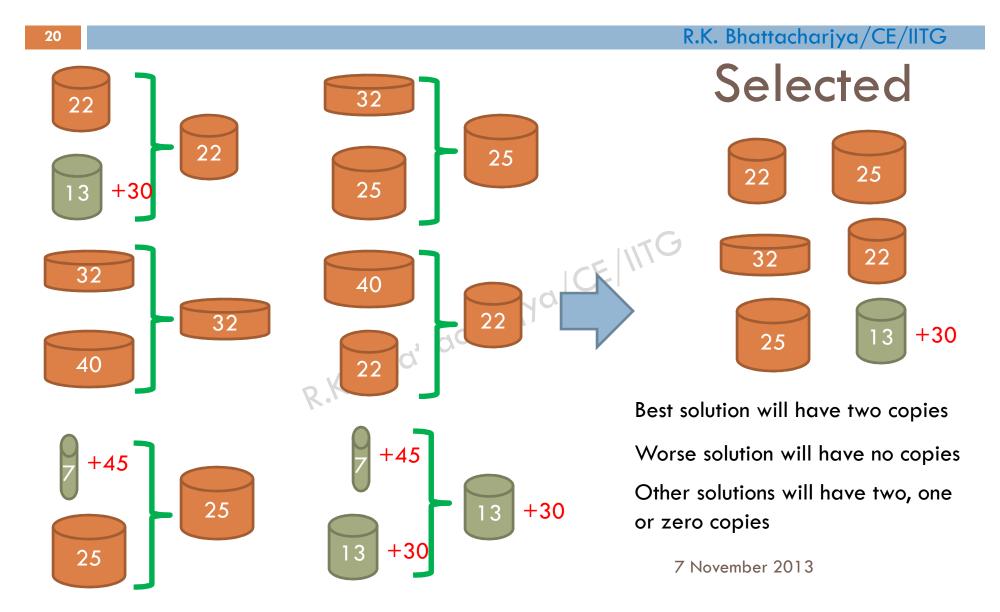
- Tournament selection and celling
 Roulette what 's
- Proportionate selection
- Rank selection
- Steady state selection, etc

Tournament selection

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- In tournament selection several tournaments are played among a few individuals. The individuals are chosen at random from the population.
- The winner of each tournament is selected for next generation.
- Selection pressure can be adjusted by changing the tournament size.
- Weak individuals have a smaller chance to be selected if tournament size is large.

Tournament selection

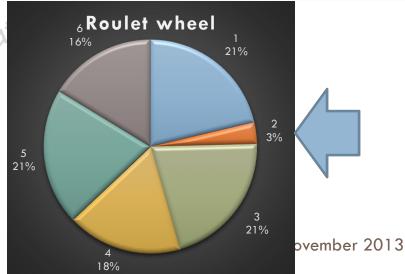


Roulette wheel and proportionate selection

Parents are selected according to their fitness values

The better chromosomes have more chances to be selected

| Chrom # | Fitness | % of RW | EC | AC |
|---------|---------|---------|------|----|
| 1 | 50 | 26.88 | 1.61 | 2 |
| 2 | 6 | 3.47 | 0.19 | 0 |
| 3 | 36 | 20.81 | 1.16 | 1 |
| 4 | 30 | 17.34 | 0.97 | 1 |
| 5 | 36 | 20.81 | 1.16 | 1 |
| 6 | 28 | 16.18 | 0.90 | 1 |
| | 186 | 100.00 | 6 | 6 |

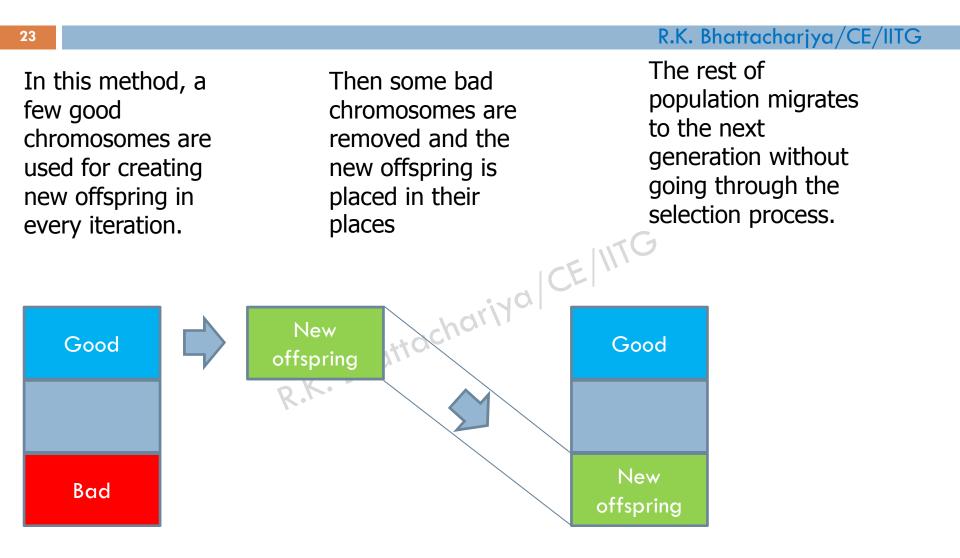


Rank selection

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| Chrom # | Fitness | Sort | Chrom # | Fitness | | Chro | om # Rank | < |
|---------|---------|---------------------|----------|----------|---------|-------|-----------|---|
| 1 | 37 | according | 1 | 37 | Assign | | 6 | |
| 2 | 6 | to fitness | 3 | 36 | raking | | 3 5 | |
| 3 | 36 | | 5 | 36 | | | 5 4 | |
| 4 | 30 | | 4 | 30 | | - | 4 3 | |
| 5 | 36 | | 6 | 28 | | | 5 2 | |
| 6 | 28 | | 2 | 6 | | | 2 1 | |
| | | | ++ achui | <i>b</i> | | | | |
| Chrom # | % of RW | Roulette whe | el | | Chrom # | EC | AC | |
| 1 | 29 | RK | | | 1 | 1.714 | 2 | |
| 3 | 24 | | | | 3 | 1.429 | 1 | |
| 5 | 19 | | | | 5 | 1.143 | 1 | |
| 4 | 14 | | | | 4 | 0.857 | 1 | |
| 6 | 10 | | | | 6 | 0.571 | 1 | |
| 2 | 5 | • 6 • 5 • 4 • 3 • 2 | =] | | 2 | 0.286 | 0 | |

Steady state selection



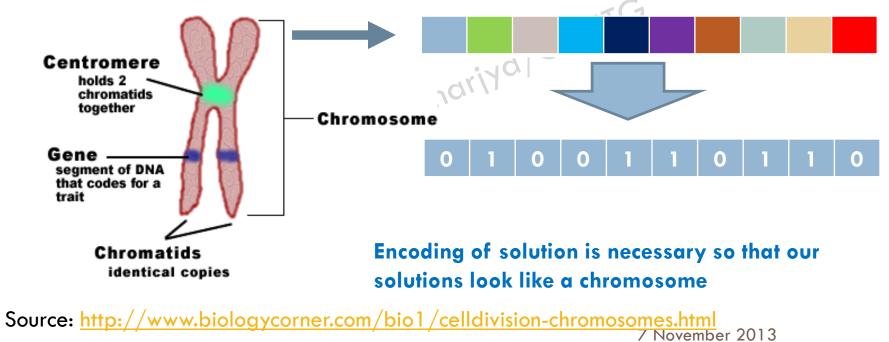
How to implement crossover

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The crossover operator is used to create new solutions from the existing solutions available in the mating pool after applying selection operator.

This operator exchanges the gene information between the solutions in the mating pool.



Encoding

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The process of representing a solution in the form of a string that conveys the necessary information.

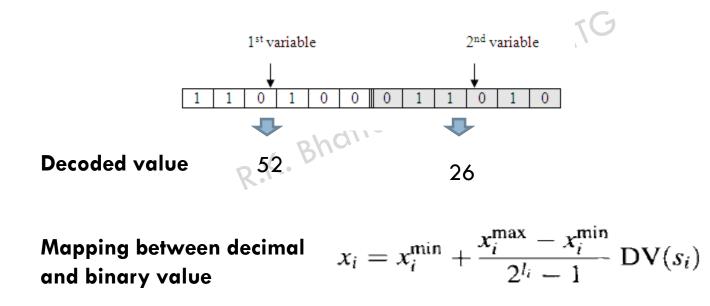
Just as in a chromosome, each gene controls a particular characteristic of the individual, similarly, each bit in the string represents a characteristic of the solution.

Encoding Methods

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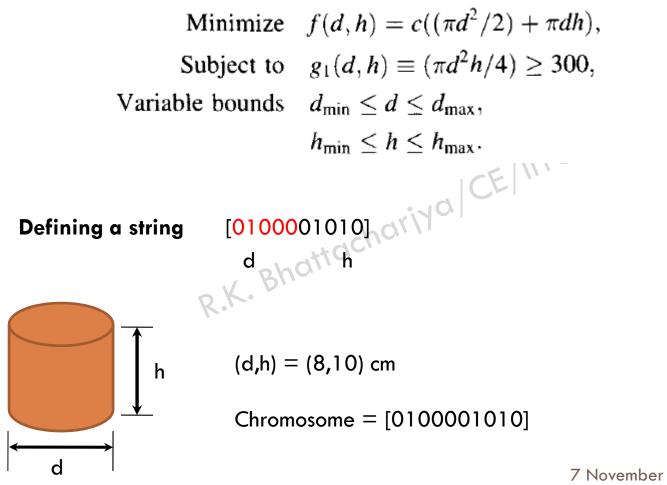
Most common method of encoding is binary coded. Chromosomes are strings of 1 and 0 and each position in the chromosome represents a particular characteristic of the problem



Encoding Methods

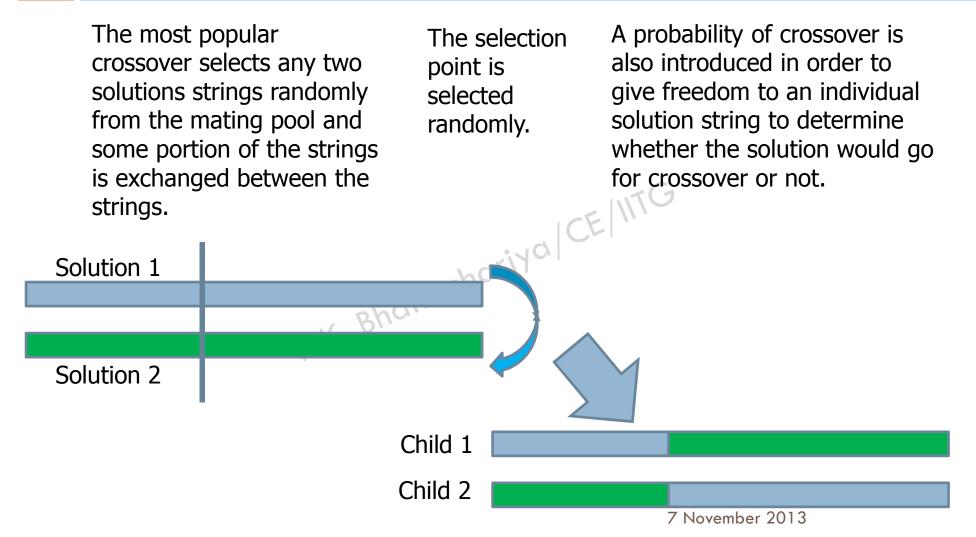
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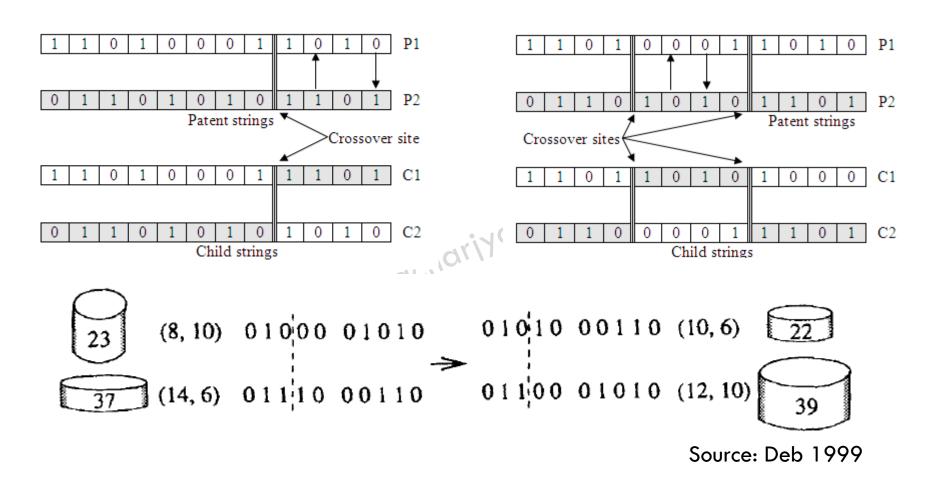
Crossover operator

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Binary Crossover

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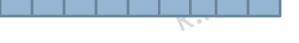
Mutation operator

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Mutation is the occasional introduction of new features in to the solution strings of the population pool to maintain diversity in the population.

Though crossover has the main responsibility to search for the optimal solution, mutation is also used for this purpose. Phattachariya





Before mutation

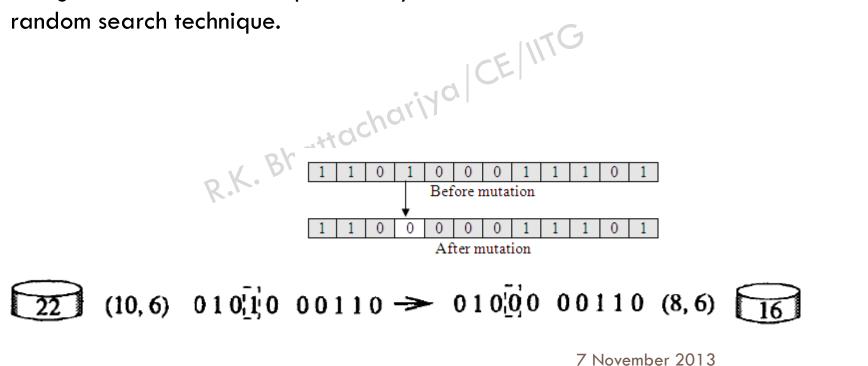
After mutation

Binary Mutation

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Source: Deb 1999

- Mutation operator changes a 1 to 0 or vise versa, with a mutation probability of .
- The mutation probability is generally kept low for steady convergence.
- A high value of mutation probability would search here and there like a random search technique.



Elitism

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Crossover and mutation may destroy the best solution of the population pool

Elitism is the preservation of few best solutions of the population pool

Elitism is defined in percentage or in number

Nature to Computer Mapping

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| Nature | Computer |
|------------|-------------------------------|
| Population | Set of solutions |
| Individual | Solution to a problem |
| Fitness | Quality of a solution |
| Chromosome | Encoding for a solution |
| Gene | Part of the encoding solution |
| R.N. | |

An example problem

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Maximize $f(x) = \sin(x)$

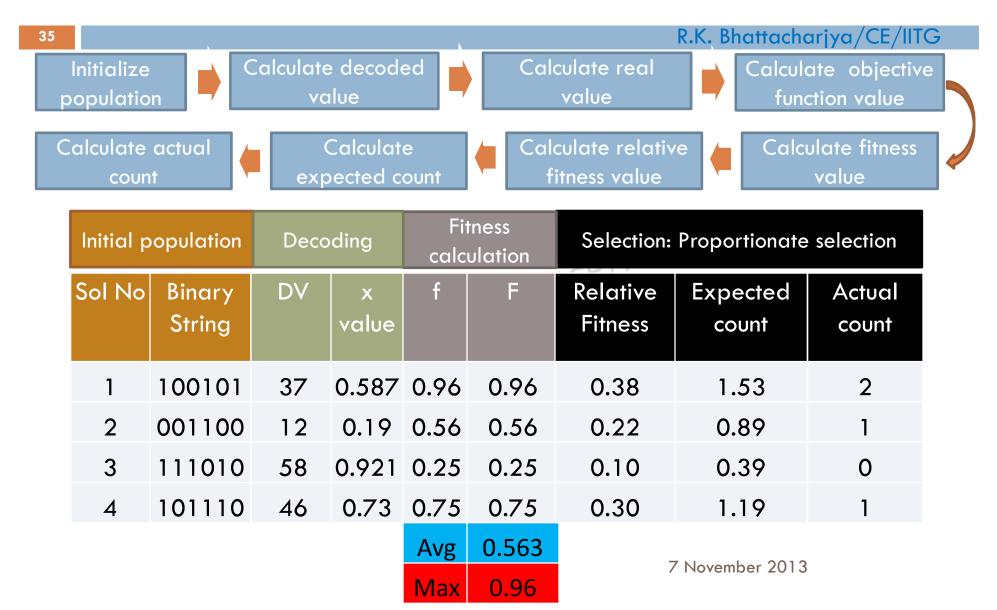
 $0 \leq x \leq \pi$

Consider 6 bit string to represent the solution, then JI 4 Itachariya CE ITG 000000 = 0 and $111111 = \pi$

```
Assume population size of 4
```

Let us solve this problem by hand calculation

An example problem



An example problem: Crossover

| | | | | | | R.K. | Bhattachc |
|-----------|-----------------|---------|-------------------------------------|----|------------|------|------------|
| atting po | ool | Ŗ | Candom generation crossover site | | | Ne | w populati |
| | Crosso | over: S | Single point | | | | |
| Sol No | Matting pool | CS | New Binary String | DV | x value | f | F |
| 1 | 100101 | 3 | 100100 | 36 | 0.57 | 0.97 | 0.97 |
| 2 | 001100 | 3 | 001101 | 13 | 0.21 | 0.60 | 0.60 |
| 3 | 100101 | 2 | 101110 | 46 | 0.73 | 0.75 | 0.75 |
| 4 | 101110 | 2 | 100101 | 37 | 0.59 | 0.96 | 0.96 |
| | | | | | | Avg | 0.8223 |
| | | | | | | Max | 0.97 |

An example problem: Mutation

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| | Muto | ation | | | | |
|--------|-------------------------------|------------------------------|----|------------|------|-------|
| Sol No | Population after crossover | Population after mutation | DV | x value | f | F |
| 1 | 100100 | 100000 | 32 | 0.51 | 1.00 | 1.00 |
| 2 | <mark>0</mark> 01101 | 101101 | 45 | 0.71 | 0.78 | 0.78 |
| 3 | 10 <mark>1</mark> 110 | 100110 | 38 | 0.60 | 0.95 | 0.95 |
| 4 | 100101 | 101101 | 45 | 0.71 | 0.78 | 0.78 |
| | 9 | | | | Avg | 0.878 |
| | | | | | Max | 1.00 |

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Disadvantage of binary coded GA

- more computation
- Iower accuracy

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- Ionger computing time
- riya CE ITG solution space discontinuity
- hamming cliff

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The standard genetic algorithms has the following steps

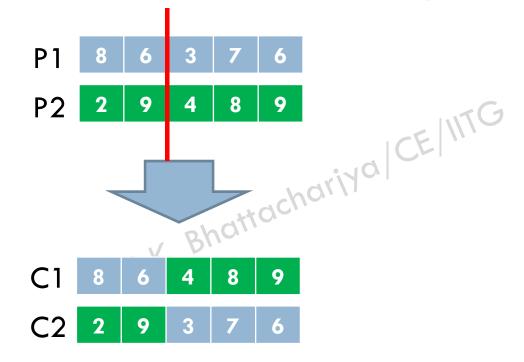
- Choose initial population 1.
- Assign a fitness function 2.
- Perform elitism 3.

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- 4.
- 5.
- Perform mutation to chariya 6. R.K. Bha
- In case of standard Genetic Algorithms, steps 5 and 6 require bitwise manipulation.

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Simple crossover: similar to binary crossover



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Linear Crossover

- Parents: $(x_1, ..., x_n)$ and $(y_1, ..., y_n)$
- Select a single gene (*k*) at random
- Three children are created as,

$$(x_1, ..., x_k, 0.5 \cdot y_k + 0.5 \cdot x_k, ..., x_n)$$

$$(x_1, ..., x_k, 1.5 \cdot y_k - 0.5 \cdot x_k, ..., x_n)$$

$$(x_1, ..., x_k, -0.5 \cdot y_k + 1.5 \cdot x_k, ..., x_n)$$

• From the three children, best two are selected for the next generation

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Single arithmetic crossover

- Parents: $(x_1, ..., x_n)$ and $(y_1, ..., y_n)$ ullet
- Select a single gene (k) at random ullet
- child₁ is created as, ullet

child₁ is created as,

$$(x_1, ..., x_k, \boldsymbol{\alpha} \cdot y_k + (1 - \boldsymbol{\alpha}) \cdot x_k, ..., x_n)$$

reverse for other child. e.g. with $\alpha = 0.5$ ullet

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Simple arithmetic crossover

- Parents: $(x_1, ..., x_n)$ and $(y_1, ..., y_n)$
- Pick random gene (k) after this point mix values
- child₁ is created as: $(x_1, ..., x_k, \alpha \cdot y_{k+1} + (1-\alpha) \cdot x_{k+1}, ..., \alpha \cdot y_n + (1-\alpha) \cdot x_n)$
- reverse for other child. e.g. with $\alpha = 0.5$

0.7

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Whole arithmetic crossover

- Most commonly used •
- Parents: (x_1, \dots, x_n) and (y_1, \dots, y_n) child₁ is: •
- •

$$\alpha \cdot \overline{x} + (1 - \alpha) \cdot \overline{y}$$

reverse for other child. e.g. with $\alpha = 0.5$ ullet



Simulated binary crossover

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Developed by Deb and Agrawal, 1995)

$$\begin{split} x_i^{(1,t+1)} &= 0.5 \left[\left(1 + \beta_{q_i} \right) x_i^{(1,t)} + \left(1 - \beta_{q_i} \right) x_i^{(2,t)} \right] \\ x_i^{(2,t+1)} &= 0.5 \left[\left(1 - \beta_{q_i} \right) x_i^{(1,t)} + \left(1 + \beta_{q_i} \right) x_i^{(2,t)} \right] \\ \beta_{q_i} &= \begin{cases} (2u_i)^{\frac{1}{n_c+1}} &, & \text{if } u_i \leq 0.5 \\ \left(\frac{1}{2(1-u_i)} \right)^{\frac{1}{n_c+1}}, & \text{otherwise} \end{cases} \end{split}$$

Where, u_i a random number

 n_c is a parameter that controls the crossover process. A high value of the parameter will create near-parent solution

Random mutation

 $y_i^{(1,t+1)} = u_i (x_i^u - x_i^l)$

Where u_i is a random number between [0,1]

$$y_i^{(1,t+1)} = x_i^{1,t+1} + (u_i - 0.5)\Delta_i \text{ ariya} |CE| |MG|$$

R.K. Bhottoolian Arian is the user defined maximum

Where, Δ_i is the user defined maximum perturbation

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Normally distributed mutation

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A simple and popular method

$$y_i^{(1,t+1)} = x_i^{1,t+1} + N(0,\sigma_i)$$

Where $N(0, \sigma_i)$ is the Gaussian probability distribution with zero mean

Polynomial mutation

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Deb and Goyal, 1996 proposed

$$y_i^{1,t+1} = x_i^{1,t+1} + (x_i^u - x_i^l)\delta_i$$

$$\delta_i = \begin{bmatrix} (2u_i)^{1/(\eta_m+1)} - 1 & \text{If } u_i < 0.5 \\ 1 - (2(1-u_i))^{1/(\eta_m+1)} & \text{If } u_i \ge 0.5 \end{bmatrix}$$

Multi-modal optimization

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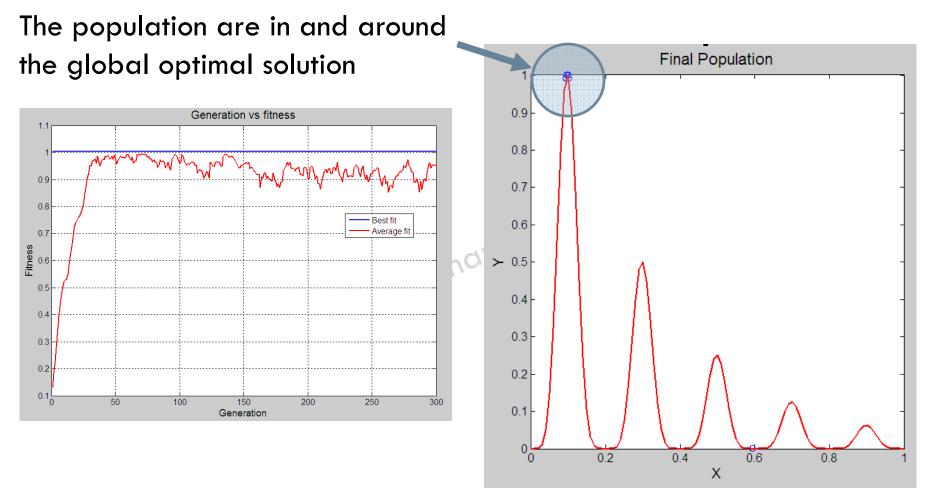
Minimize
$$f(x) = 2^{-2(\frac{(x-0.1)}{0.8})^2 \sin^6(5\pi x)}$$

 $0 \le x \le 1$
Solve this problem using
simple Genetic Algorithms

After Generation 200

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Multi-modal optimization

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Simple modification of Simple Genetic Algorithms can capture all the optimal solution of the problem including global optimal solutions

Basic idea is that reduce the fitness of crowded solution, which can be implemented using following three steps.

Sharing function
$$Sh(d_{ij}) = \begin{cases} 1 - (d_{ij}/\sigma), & \text{if } d_{ij} < \sigma; \\ 0, & \text{otherwise.} \end{cases}$$

Niche count $nc_i = \sum_{j=1}^{N} Sh(d_{ij})$
Modified fitness $f'_i = \frac{f_i}{nc_i}$

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Hand calculation

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 $0 \le x \le 2$

| Sol | String | Decoded value | X | f |
|-----|--------|------------------|-------|-------|
| 1 | 110100 | 52 | 1.651 | 0.890 |
| 2 | 101100 | 44 | 1.397 | 0.942 |
| 3 | 011101 | 29 | 0.921 | 0.246 |
| 4 | 001011 | 11 | 0.349 | 0.890 |
| 5 | 110000 | 48 | 1.524 | 0.997 |
| 6 | 101110 | 46 | 1.460 | 0.992 |

Distance table

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| dij | 1 | 2 | 3 | 4 | 5 | 6 |
|-----|-------|-------|-------|-------|-------|-------|
| 1 | 0 | 0.254 | 0.73 | 1.302 | 0.127 | 0.191 |
| 2 | 0.254 | 0 | 0.476 | 1.048 | 0.127 | 0.063 |
| 3 | 0.73 | 0.476 | 0 | 0.572 | 0.603 | 0.539 |
| 4 | 1.302 | 1.048 | 0.572 | 0 | 1.175 | 1.111 |
| 5 | 0.127 | 0.127 | 0.603 | 1.175 | 0 | 0.064 |
| 6 | 0.191 | 0.063 | 0.539 | 1.111 | 0.064 | 0 |

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Sharing function values

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| sh(dij) | 1 | 2 | 3 | 4 | 5 | 6 | nc |
|---------|-------|-------|-------|---|-------|-------|-------|
| 1 | 1 | 0.492 | 0 | 0 | 0.746 | 0.618 | 2.856 |
| 2 | 0.492 | 1 | 0.048 | 0 | 0.746 | 0.874 | 3.16 |
| 3 | 0 | 0.048 | 1 | 0 | 0 | 0 | 1.048 |
| 4 | 0 | 0 | 0 | 1 | 0 | 0 | 1 |
| 5 | 0.746 | 0.746 | 0 | 0 | 1 | 0.872 | 3.364 |
| 6 | 0.618 | 0.874 | 0 | 0 | 0.872 | 1 | 3.364 |

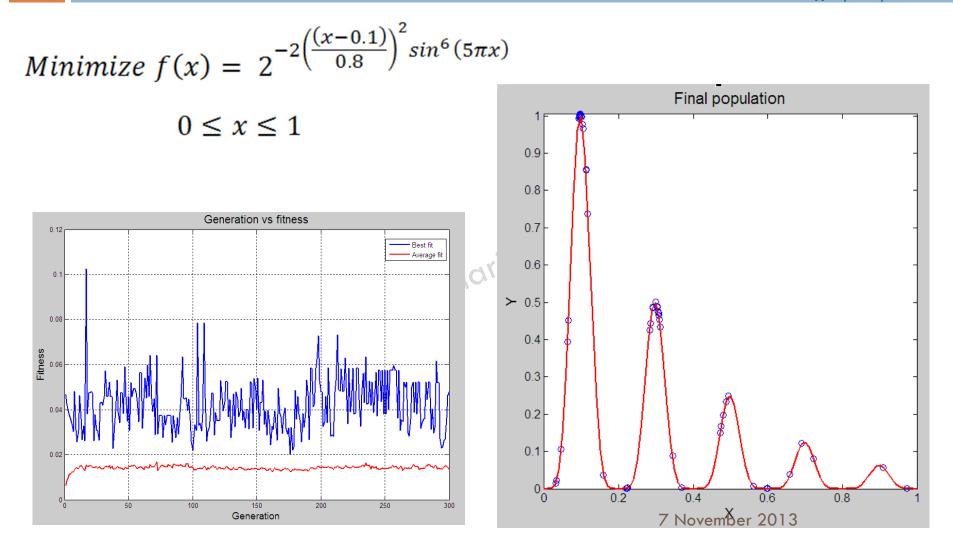
Sharing fitness value

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| Sol | String | Decoded value | x | f | nc | f' |
|-----|--------|------------------|-------|-------|-------|-------|
| 1 | 110100 | 52 | 1.651 | 0.890 | 2.856 | 0.312 |
| 2 | 101100 | 44 | 1.397 | 0.942 | 3.160 | 0.300 |
| 3 | 011101 | 29 | 0.921 | 0.246 | 1.048 | 0.235 |
| 4 | 001011 | 11 | 0.349 | 0.890 | 1.000 | 0.890 |
| 5 | 110000 | 48 | 1.524 | 0.997 | 3.364 | 0.296 |
| 6 | 101110 | 46 | 1.460 | 0.992 | 3.364 | 0.295 |

Solutions obtained using modified fitness value

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Evolutionary Strategies

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ES use real parameter value

□ ES does not use crossover operator

It is just like a real coded genetic algorithms with selection and mutation operators only

CEINTG

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In each iteration one parent is used to create one offspring by using Gaussian mutation operator

- \Box Step 1: Choose a initial solution x and a mutation strength σ
- Step2: Create a mutate solution

$$y = x + N(0, \sigma)$$

- $y = x + N(0, \sigma)$ $\Box \text{ Step 3: If } f(y) < f(x), \text{ replace } x \text{ with } y$
- □ Step4: If termination criteria is satisfied, stop, else go to step 2

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 \square Strength of the algorithm is the proper value of σ

- Rechenberg postulate
 - The ratio of successful mutations to all the mutations should be 1/5. If this ratio is greater than 1/5, increase mutation strength. If it is less than 1/5, decrease the mutation strength.

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- A mutation is defined as successful if the mutated offspring is better than the parent solution.
- □ If P_s is the ratio of successful mutation over *n* trial, Schwefel (1981) suggested a factor $C_d = 0.817$ in the following σ update rule

$$\sigma^{t+1} = \begin{cases} C_d \sigma^t & \text{if } P_s < 1/5 \\ \frac{1}{C_d} \sigma^t & \text{if } P_s < 1/5 \\ \sigma^t & \text{if } P_s = 1/5 \end{cases}$$

Matlab code

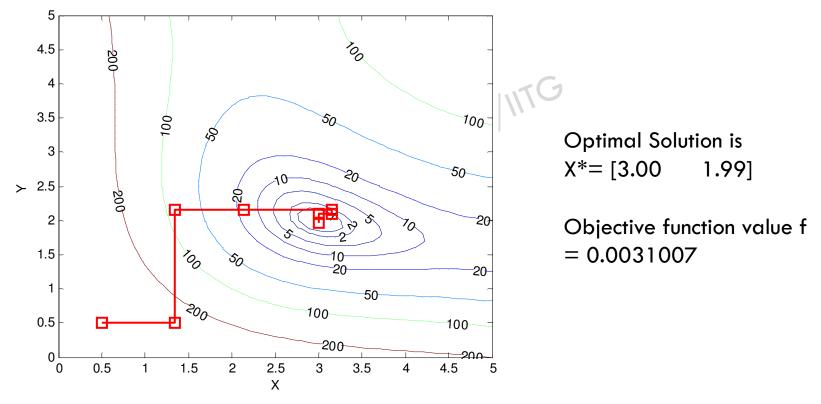
```
sigma = 1;
                                  function [f] = objfunc( x )
 x0 = [1 1];
                                  f = (x(1)^{2}+x(2)-11)^{2}+(x(1)+x(2)^{2}-7)^{2};
 [n m] = size(x0);
                                  end
for j=1:1000
 for i = 1:m
     fO = objfunc(xO);
     x1 = x0;
     x1(i) =x0(i) *randn(1) *sigma;
     f1 = objfunc(x1);
      if (f1<f0)
          x0 = x1;
     end
-end
 end
 disp(['Optimal solution X= ', num2str(x0)]);
                                                 / November 2013
```

```
% This programme will implement 1+1 ES
bx = [0 5];  Upper bound
bv = [0 5];  Lower bound
plotfunction (bx, by) % Ploting the function between upper bound and lower
bound defined above
hold on:
x0 = [0.5 0.5]; % Starting point or initial solution
sigma = 5; % Define sigma value
imax = 3000; % maximum iteration
k =0: % An counter
success =0: % Success counter
[n m] = size(x0);
x11 =x0; % x11 will store solution of all the iteration
for j=1:imax % The program will terminate after 3000 iteration
   k=k+1:
for i =1:m
   f0 = objfunc(x0); % objjunc will calculate the objective function value
   x1 = x0;
   x1(i) =x0(i)*randn(1)*sigma; % Will generate a new solution
   f1 = objfunc(x1);
   if (f1<f0)
       x0 = x1:
       success = success+1;
   end
   x11 = [x11; x0];
end
% Updating sigma value as per Rechenberg postulate after every 20 iterations
if(k==20)
    if(success/k>1/5)
        sigma = sigma/0.817;
   else
        sigma = sigma*0.817;
   end
   k=0;
   success =0;
end
end
plot(x11(:,1), x11(:,2),'-rs','linewidth',2,'MarkerSize',10); % plot the
solution
                                                        7 November 2013
disp(['Optimal solution X= ', num2str(x0)]);
disp(['Optimal function value f= ', num2str(f0)]);
```

Some results of 1+1 ES

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Minimize
$$f = (x_1^2 + x_2 - 11)^2 + (x_1 + x_2^2 - 7)^2$$



Multimember ES

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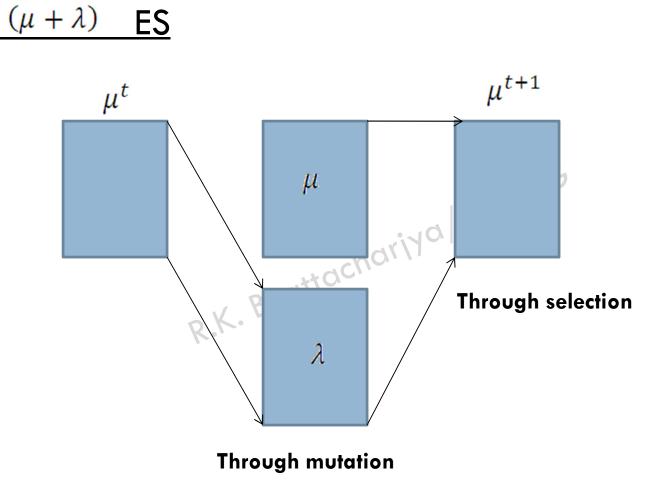
$(\mu + \lambda)$ ES

Step 1: Choose an initial population of μ solutions and mutation strength σ Step 2: Create λ mutated solution $y^i = x^i + N(0, \sigma)$ Step 3: Combine x and y, and choose the best solutions μ

Step4: Terminate? Else go to step 2

Multimember ES

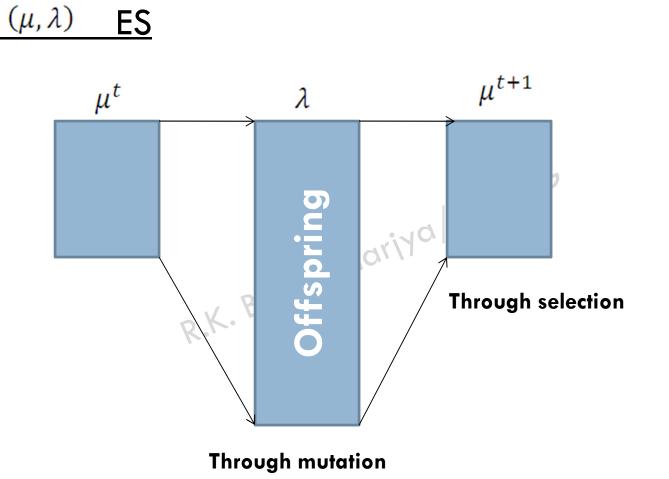
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Multimember ES

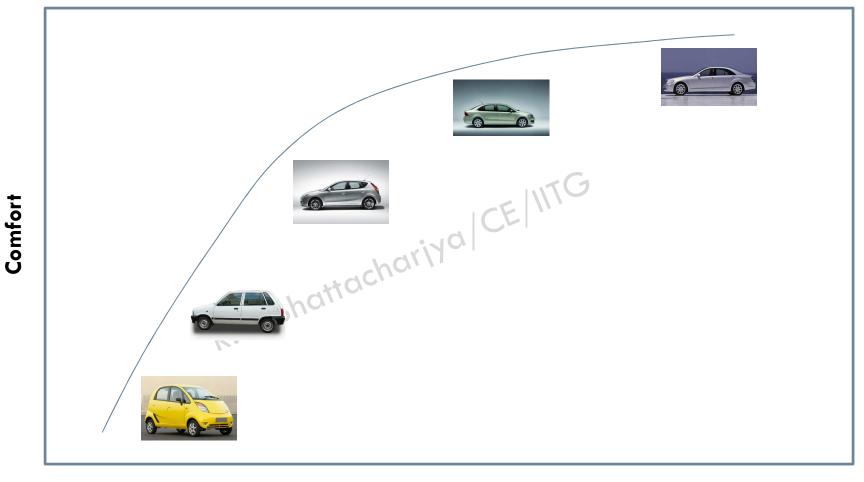
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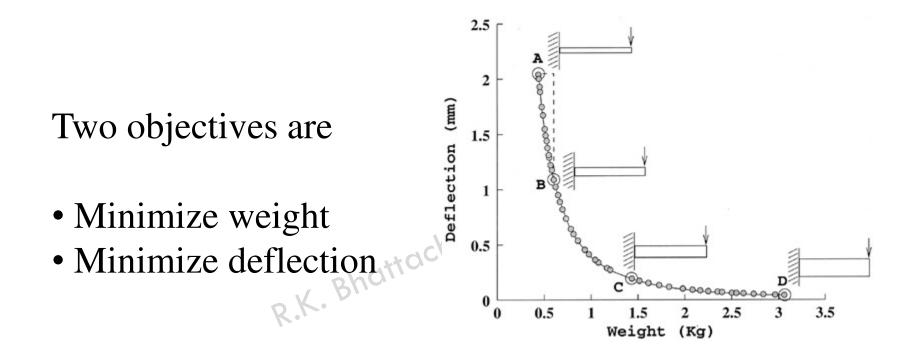
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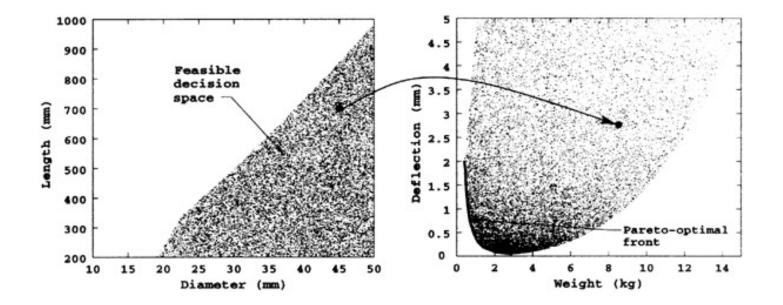
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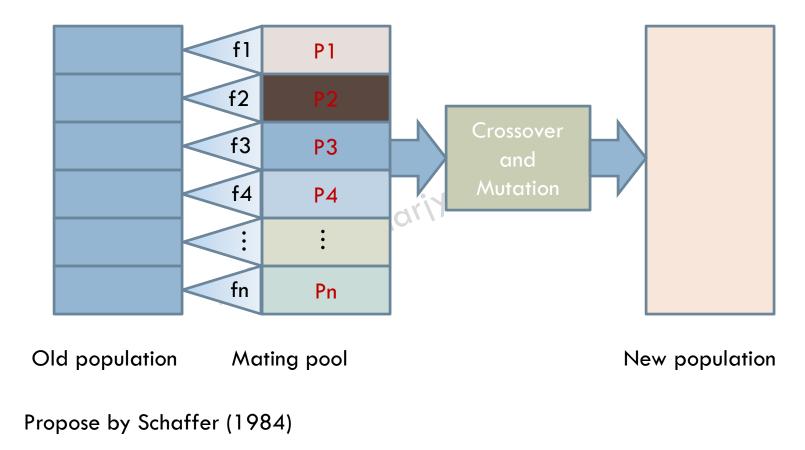
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- More than one objectives
- Objectives are conflicting in nature
- Dealing with two search space
 - Decision variable space
 - Objective space
- ICEINTG Unique mapping between the objectives and often the mapping is non-linear
- Properties of the two search space are not similar
- Proximity of two solutions in one search space does not mean a proximity in other search space

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Vector Evaluated Genetic Algorithm (VEGA)



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Non-dominated selection heuristic

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Give more emphasize on the non-dominated solutions of the population

This can be implemented by subtracting \in from the dominated solution fitness value

Suppose N' is the number of sub-population and n' is the non-dominated solutions. Then total reduction is $(N' - n') \in$.

The total reduction is then redistributed among the non-dominated solution by adding an amount $(N' - n') \in /n'$

This method has two main implications

Non-dominated solutions are given more importance

Additional equal emphasis has been given to all the non-dominated solution

Weighted based genetic algorithm (WBGA)

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The fitness is calculated

$$F = \sum_{j=1}^{M} w_j \frac{f_i - f_j^{min}}{f_j^{max} - f_j^{min}}$$

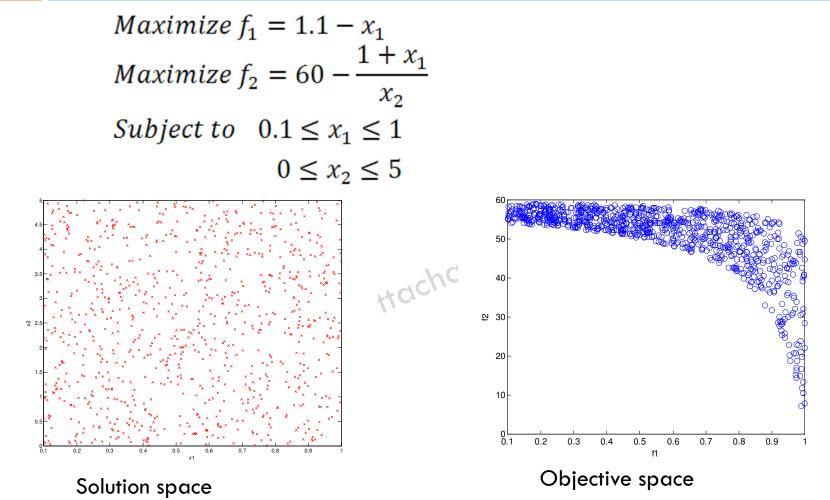
The spread is maintained using the sharing function approach

Sharing function
$$Sh(d_{ij}) = \begin{cases} 1 - (d_{ij}/\sigma), & \text{if } d_{ij} < \sigma \\ 0, & \text{otherwise.} \end{cases}$$

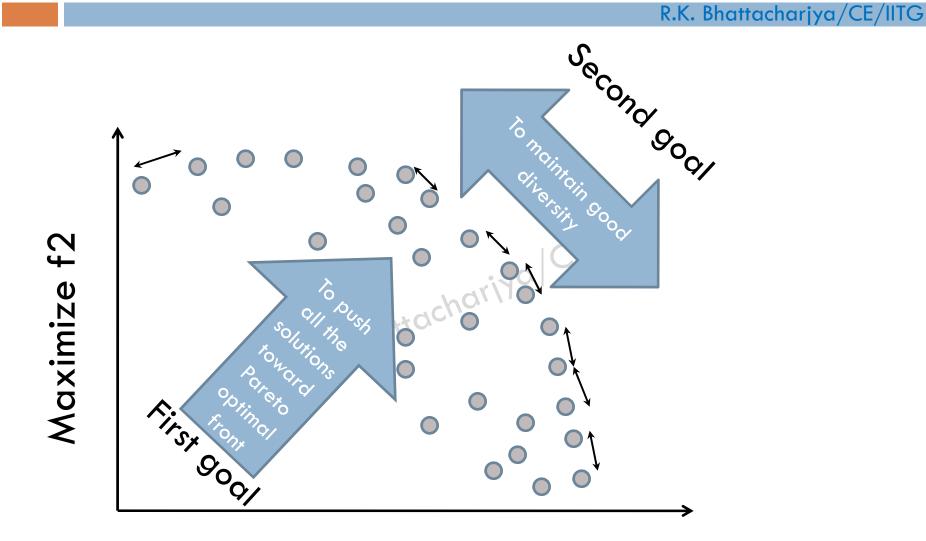
 $nc_i = \sum_{i=1}^{n} Sh(d_{ij})$

Modified fitness $F' = \frac{F}{nc}$

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Maximize f1

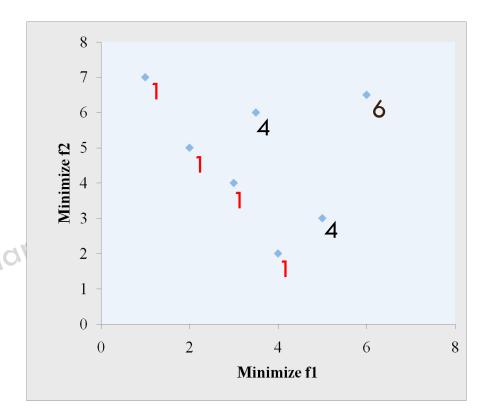
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Fonseca and Fleming (1993) first introduced multiple objective genetic algorithm (MOGA)

The assigned fitness value based on the non-dominated ranking.

The rank is assigned as $r_i = 1 + n_i$ where r_i is the ranking of the i^{th} solution and n_i is the number of solutions that dominate the solution.



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- Fonseca and Fleming (1993) maintain the diversity among the non-dominated solution using niching among the solution of same rank.
- The normalize distance was calculated as,

$$d_{i,j} = \sqrt{\sum_{k=1}^{M} \left(\frac{f_k^i + f_k^j}{f_k^{max} - f_k^{min}}\right)^2}$$

The niche count was calculated as,

$$nc_i = \sum_{j=1}^{\mu(r_i)} Sh(d_{ij})$$

NSGA

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- Srinivas and Deb (1994) proposed NSGA
- The algorithm is based on the non-dominated sorting.

CENTG

The spread on the Pareto optimal front is maintained using sharing function

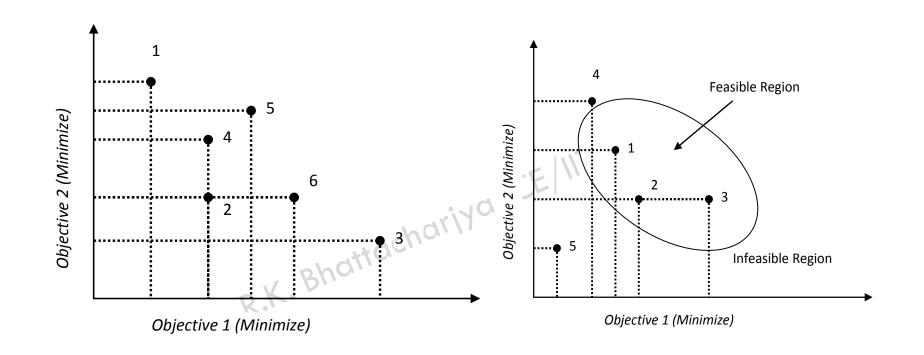
$$d_{i,j} = \sqrt{\sum_{k=1}^{P_1} \left(\frac{x_k^i - x_k^j}{x_k^{max} - x_k^{min}} \right)^2}$$

NSGA II

- Non-dominated Sorting Genetic Algorithms
 - NSGA II is an elitist non-dominated sorting Genetic Algorithm to solve multi-objective optimization problem developed by Prof. K. Deb and his student at IIT Kanpur.
 - It has been reported that NSGA II can converge to the global Pareto-optimal front and can maintain the diversity of population on the Pareto-optimal front

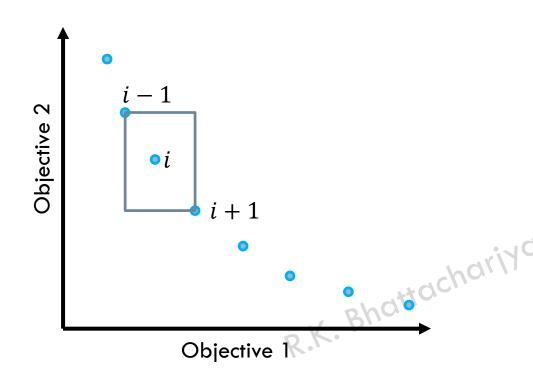
Non-dominated sorting

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Calculation crowding distance

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Cd, the crowded distance is the perimeter of the rectangle constituted by the two neighboring solutions

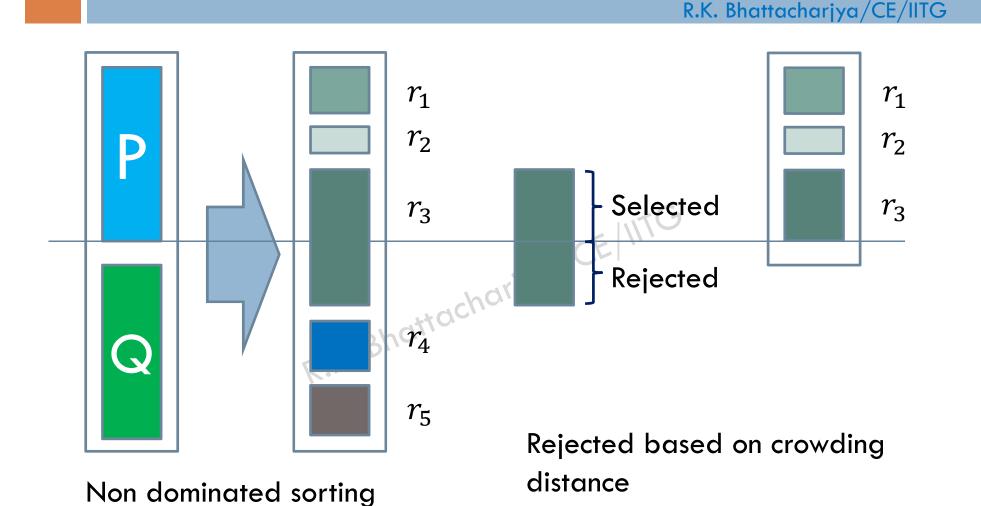
Cd value more means that the solution is less crowded

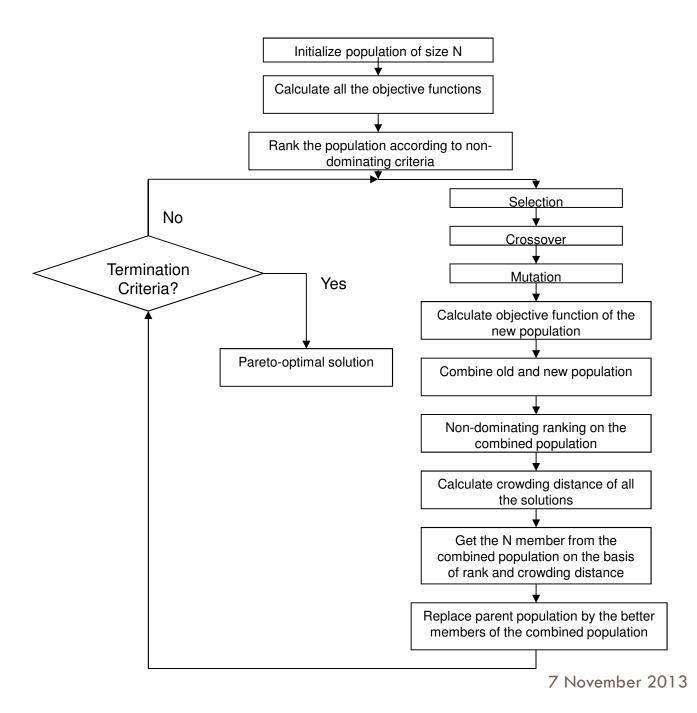
Cd value less means that the solution is more crowded

Crowded tournament operator

- □ A solution i wins a tournament with another solution j_i
 - If the solution i has better rank than j, i.e. $r_i < r_j$
 - If they have the same rank, but i has a better crowding distance than j, i.e. $r_i = r_j$ and $d_i > d_j$.

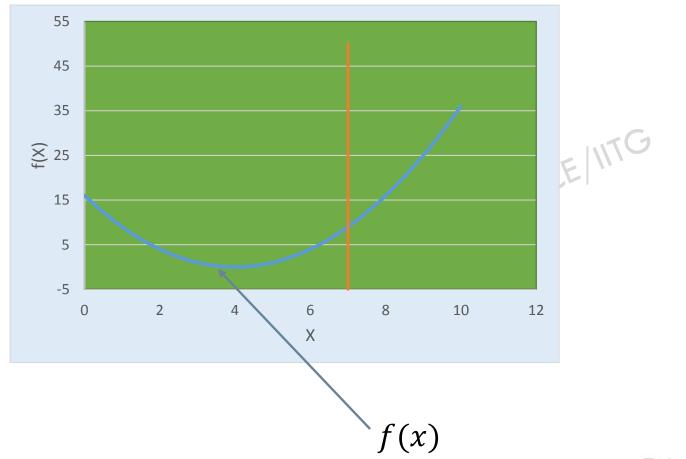
Replacement scheme of NSGA II





Constraints handling in GA

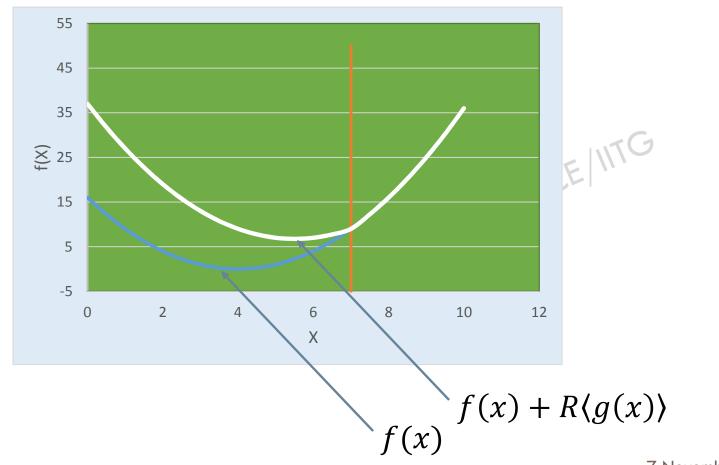
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Constraints handling in GA

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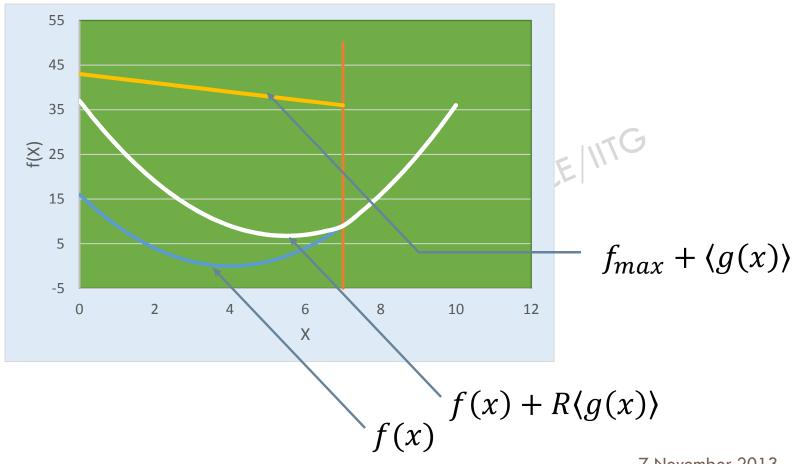
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Constraints handling in GA

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Constrain handling in GA

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Minimize f(X)Subject to $g_{j}(x) \leq 0 \qquad j = 1,2,3,...,J$ $h_{k}(x) = 0 \qquad k = 1,2,3,...,K$ o's approach $F = f(X)_{R}$ Deb's approach If X is feasible $= f_{max} + \sum_{j=1}^{J} \langle g_j(x) \rangle + \sum_{k=1}^{K} |h_k(x)|$ Otherwise

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