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Problem: [Arora, Barak - Exercise 2.15] In the VERTEX COVER problem, we are given an undirected graph G and an integer K and have to decide whether there is a subset S of at most K vertices such that every edge (i,j) of G , at least one of i or j is in S (such a subset is called a *vertex cover* of G). Prove that this problem is NP-complete.

Solution :

We know that the following problem is NP-complete :

P1: An independent set of a graph $G = (V, E)$ is a $V_I \subseteq V$ such that no two vertices in V_I share an edge. Does G have an independent set of size at least k ?

We will reduce it to following :

P2: A vertex cover of a graph $G = (V, E)$ is a $V_C \subseteq V$ such that every $(a, b) \in E$ is incident to at least a $u \in V_C$. Does G have a vertex cover of size at most k ?

A : Vertex Cover is in NP.

Given V_C , vertex cover of $G = (V, E)$, $|V_C| \leq k$

We can check in $O(|E| + |V|)$ that V_C is a vertex cover for G . How?

For each vertex $\in V_C$, remove all incident edges.

Check if all edges were removed from G .

Thus, Vertex Cover \in NP

B. Independent Set Problem can be reduced to Vertex Cover Problem in Polynomial Time.

Given a general instance of IS: $G' = (V', E')$, k'

Construct a specific instance of VC: $G = (V, E)$, k

$V = V'$

$E = E'$

$(G = G')$

$k = |V'| - k'$

This transformation is polynomial:

Constant time to construct $G = (V, E)$

$O(|V|)$ time to count the number of vertices

Proof : G has an independent set V_I of size k iff VC has a vertex cover V_C of size at most k' .

Consider two sets I and J s.t. $I \cap J = \emptyset$ and $I \cup J = V = V'$

Given any edge (u, v) , one of the following four cases holds:

$u, v \in I$

$u \in I$ and $v \in J$

$u \in J$ and $v \in I$

$u, v \in J$

Assume that I is an independent set of G then:

Case 1 cannot be : (vertices in I cannot be adjacent)

In cases 2 and 3, (u, v) has exactly one endpoint in J .

In case 4, (u, v) has both endpoints in J .

In cases 2, 3 and 4, (u, v) has at least one endpoint in J .

Thus, vertices in J cover all edges of G . Also: $|I| = |V| - |J|$ since $I \cap J = \emptyset$ and $I \cup J = V = V$

Thus, if I is an independent set of G , then J is a vertex cover of $G (= G)$.

Similarly, we can prove that if J is a vertex cover for G , then I is an independent set for G .

Hence Proved.