

Discrete Mathematics

Benny George K

Department of Computer Science and Engineering

Indian Institute of Technology Guwahati

`ben@iitg.ernet.in`

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Elementary Concepts

Let A and B be sets. Let $A_i, i \in \mathcal{I}$ be an indexed family of sets, i.e. for each $i \in \mathcal{I}$ we have sets A_i (Assume $\mathcal{I} \neq \emptyset$)

► Union

$$A \cup B \triangleq \{x \mid x \in A \text{ or } x \in B\}$$

$$\bigcup_{i \in \mathcal{I}} A_i \triangleq \{x \mid \exists j, x \in A_j\}$$

► Intersection

$$A \cap B \triangleq \{x \mid x \in A \text{ and } x \in B\}$$

$$\bigcap_{i \in \mathcal{I}} A_i \triangleq \{x \mid \forall j, x \in A_j\}$$

- Set Difference (Also written as $A - B$ and called as relative complement of B relative to A and shortened as B^c when A is clear from the context)

$$A \setminus B \triangleq \{x \mid x \in A \text{ and } x \notin B\}$$



Elementary Concepts

- ▶ Symmetric Difference

$$A \oplus B \triangleq (A \setminus B) \cup (B \setminus A)$$

- ▶ Power set of a set S (Written as $\mathcal{P}(S)$ or 2^S)

$$\mathcal{P}(S) \triangleq \{x \mid x \subseteq S\}$$

- ▶ DeMorgan's rule.

$$(A \cup B)^c = A^c \cap B^c$$

$$(A \cap B)^c = A^c \cup B^c$$



Finite and infinite set.

- ▶ For any given set A , define A^+ called **successor** of A as below.

$$A^+ \triangleq A \cup \{A\}$$

- ▶ We can start with the empty set \emptyset , repeatedly apply the successor operation and construct a sequence of sets.
- ▶ The first few sets in this sequence will be $\emptyset, \{\emptyset\}, \{\emptyset, \{\emptyset\}\}, \{\emptyset, \{\emptyset\}, \{\emptyset, \{\emptyset\}\}} \dots$
- ▶ We shall name these sets as $0, 1, 2, 3, \dots$
- ▶ Let us now define the set \mathbb{N} to be the set which
 - ▶ contains 0.
 - ▶ Whenever it contains the element A , it contains A^+ as well.
- ▶ \mathbb{N} constructed as above is an “infinite” set but we will formally define that term in the next page.



Finite and Infinite Sets

- ▶ A set S is said to be **finite** if there is a bijection (one to one correspondence) between S and an element of \mathbb{N} .
- ▶ n is said to be the **cardinality** or *size* of the set S and is denoted by $|S|$.
- ▶ A set S is said to be **infinite** if it is not finite.
- ▶ A set S is said to be **countably infinite** if there exists a bijection between S and \mathbb{N} .
- ▶ A set S is said to be **countable** or **enumerable** if it is finite or countably infinite. (e.g. \mathbb{Z}, \mathbb{Q}).
- ▶ A set is said to be **uncountable** if it is not countable. (e.g. $\mathbb{R}, [0, 1]$, set of irrationals etc.)
- ▶ We say that two sets S_1 and S_2 are of same cardinality if there is a bijection from S_1 to S_2 .



Cardinality of sets

Integers(\mathbb{Z}) form a countable set

Consider the map f from \mathbb{Z} to \mathbb{N} given by $f(x) = 2x$ if $x \geq 0$ and $f(x) = 2x - 1$ if $x < 0$.

Rationals form a countable set.

Every positive rational number is of the form $p/q, q \neq 0$. List the rationals in increasing order of $p + q$. (We can do this because there are only finitely many positive integral solutions for the equation $p + q = k$ for any fixed k). Negative rationals can be similarly enumerated then can combine these enumerations as we did for integers.

Finite subsets of \mathbb{N} is a countable set.

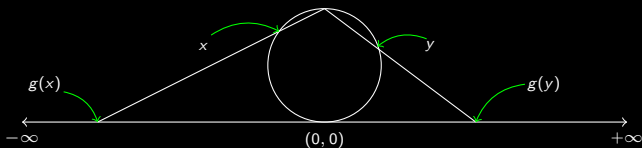
Enumerate in the increasing order of sum of elements in the subset.



Cardinality of \mathbb{R}

Power set of \mathbb{N} is of the same cardinality as \mathbb{R}

- ▶ We need to exhibit a bijection from $\mathcal{P}(\mathbb{N})$ to \mathbb{R} .
- ▶ We shall first exhibit a bijection f from $\mathcal{P}(\mathbb{N})$ to $[0, 1]$ and then a bijection g from $[0, 1]$ to \mathbb{R} . ($g \circ f$) will then be a bijection from $\mathcal{P}(\mathbb{N})$ to \mathbb{R} .
- ▶ Verify that $f(S) = \sum_{s \in S} 2^{-s}$ is a bijection from $\mathcal{P}(\mathbb{N})$ to \mathbb{R}
- ▶ The diagram below shows a bijection between $[0, 1]$ and \mathbb{R} . The circle in the diagram is the $[0, 1]$ interval rolled into a circle.



Cantor's theorem

Theorem

Cardinality of \mathbb{N} is not the same as the cardinality of \mathbb{R}

Proof

(We will show that there are no bijections from \mathbb{N} to $[0, 1]$).

- ▶ For contradiction assume that f is a bijection from \mathbb{N} to $[0, 1]$.
- ▶ Let for $n \in \mathbb{N}$, $f(n) = 0.a_{n,1}a_{n,2}a_{n,3}\dots$ where $a_{n,i}$ stands for the i th digit in the decimal expansion of $f(n)$.
- ▶ Consider the number $d = d_1d_2d_3\dots$ where $d_i = a_{i,i} + 5$
- ▶ For each i , $f(i)$ differs from d in the i th digit. Thus d is not the image of any $n \in \mathbb{N}$. Thus f is not a bijection. \square
- ▶ Cardinality of \mathbb{N} is written as \aleph_0 (read as aleph not).
- ▶ Cardinality of \mathbb{R} is written as \aleph_1 (read as aleph one).



Cantor's Theorem

Theorem

For every set S , there is not bijection from S to $\mathcal{P}(S)$

Proof

- ▶ For contradiction, assume that f is bijection from S to $\mathcal{P}(S)$.
- ▶ For every $s \in S$, $f(s)$ is a subset of S . $f(s)$ may or may not contain the element s .
- ▶ Collect all the elements d from S such that the image of d does not contain d and call this set as D .
- ▶ Symbolically, $D \triangleq \{d \mid d \notin f(d)\}$
- ▶ Notice that D cannot be the image of any element $x \in S$.
 - ▶ $x \in D$ would mean $x \notin f(x) = D$.
 - ▶ $x \notin D$ would mean $x \in f(x) = D$. \square



Introduction to Propositional Logic

- ▶ Propositional Logic is a simple but useful branch of mathematical logic.
- ▶ It helps us make inferences about *propositional formulas*.
- ▶ **Propositions** are statements which has a **truth value**.
- ▶ A proposition make take either a truth value *TRUE* or a truth value *FALSE*.
- ▶ We shall denote a proposition symbolically by letters $P, Q, R \dots$
- ▶ Also we shall abbreviate *TRUTH* and *FALSE* to T and F respectively.
- ▶ From propositions using **connectives**, we form more complex statements.



Propositional Connectives

Below we give a list of commonly used propositional connectives and their meanings.

Connective	Usage	Meaning
Negation	$\neg P$	Is true if and only if P is false
Conjunction	$P \wedge Q$	Is true if and only if both P and Q are true
Disjunction	$P \vee Q$	Is false if and only if both P and Q are false
Implication	$P \Rightarrow Q$	Is false if and only if P is true and Q is false
Equivalence	$P \Leftrightarrow Q$	Is true if and only if P and Q has same truth values.

- ▶ Implication is also referred to as conditional.
- ▶ The meaning of each propositional connective can be summarized in a [truthtable](#).



Truth tables for connectives

P	Q	$\neg P$	$P \vee Q$	$P \wedge Q$	$P \Rightarrow Q$	$P \Leftrightarrow Q$
T	T	F	T	T	T	T
T	F	F	T	F	F	F
F	T	T	T	F	T	F
F	F	T	F	F	T	T

- ▶ Total number of possible connectives on two propositional symbols is $2^4 = 16$.
- ▶ Total number of possible connectives on m propositional symbols is 2^{2^m} .



Propositional Formulas and Structural Induction

Using propositional connectives and any set of propositional symbols S we can produce a lot of **formulas**.

The set of formulas \mathcal{F} consists of

- ▶ All the elements in S as well as T and F are in \mathcal{F} .
- ▶ Let φ and ψ be elements of \mathcal{F} then $(\neg\varphi)$, $(\varphi \vee \psi)$, $(\varphi \wedge \psi)$, $(\varphi \Rightarrow \psi)$ and $(\varphi \Leftrightarrow \psi)$ are also elements of \mathcal{F}

To prove theorems about set of formulae, we use principle of structural induction

Structural Induction

If $A \subseteq \mathcal{F}$ satisfies the following conditions then $A = \mathcal{F}$

- ▶ A contains all the propositional symbols as well as T and F .
- ▶ If $\alpha, \beta \in A$, then $(\neg\alpha)$, $(\alpha \vee \beta)$, $(\alpha \wedge \beta)$, $(\alpha \Rightarrow \beta)$ and $(\alpha \Leftrightarrow \beta)$ are also elements of A



Introduction

▶ Binary operation

A binary operation from a set A to a set B is a function which assigns for each ordered pair of A a unique element in B .
Mathematically this is written as below

$$f : A \times A \mapsto B$$

▶ Examples

▶ Addition: $+_{\mathbb{N}} : \mathbb{N} \times \mathbb{N} \mapsto \mathbb{N}$

▶ Addition: $+_{\mathbb{Q}} : \mathbb{Q} \times \mathbb{Q} \mapsto \mathbb{Q}$

We use the subscripts under the operation to emphasize the fact that $+$ on \mathbb{N} and \mathbb{Q} are two different functions.

▶ Minimum: $\min : \mathbb{Q} \times \mathbb{Q} \mapsto \mathbb{Q}$

▶ Multiplication, division, subtraction, exponentiation, Maximum, Concatenation . . .



Properties of operations

Let us consider a binary operations

$$\star : A \times A \mapsto B$$

- ▶ **Closure:** \star is said to be *closed* if the set B is equal to A .
- ▶ **Associativity:** Parentheses doesn't matter.

$$((a \star b) \star c) = (a \star (b \star c))$$

- ▶ **Commutativity:** Order doesn't matter. $a \star b = b \star a$
- ▶ **Existence of Identity:** A special element e for every a

$$e \star a = a \quad (\text{left identity})$$

$$a \star e = a \quad (\text{right identity})$$

If the left identity and the right identity both exist, then they must be the same (Why?) and it is called simply the *identity*.

- ▶ **Existence of Inverse** An element a^{-1} associated with each a

$$a^{-1} \star a = e \quad (\text{left inverse of } a)$$

$$a \star a^{-1} = e \quad (\text{right inverse of } a)$$



Algebraic structures

- ▶ A set A equipped with a collection of operators is called an **algebraic structure**. We shall assume that the operations are all binary operations.

Let $\mathcal{A} = (A, \star)$ be an algebraic structure.

- ▶ If \star is closed and associative then \mathcal{A} is a **semigroup**. e.g., Set of non empty strings with the concatenation operation
- ▶ If \star is closed, associative and has an identity then \mathcal{A} is a **monoid**. e.g., Set of non strings including the empty string with the concatenation operation
- ▶ If \star is closed, associative, has an identity and has inverse then \mathcal{A} is a **group**. e.g., $n \times n$ invertible matrices under matrix multiplication
- ▶ If \star is closed, associative, has an identity, has inverse and is commutative then \mathcal{A} is an **abelian group**. e.g., Integers under addition.



Generators

Let $\mathcal{A} = (A, \star)$ be an algebraic structure such that \star is closed. For $B \subseteq A$, define a set sequence of sets B_0, B_1, B_2, \dots as below

$$\begin{aligned} B_0 &\triangleq B \\ B_1 &\triangleq \{b \mid b = b_i \star b_j \text{ where } b_i, b_j \in B_0\} \cup B_0 \\ &\vdots \\ B_{i+1} &\triangleq \{b \mid b = b_i \star b_j \text{ where } b_i, b_j \in B_i\} \cup B_i \end{aligned}$$

- The set B^* defined as

$$B^* \triangleq \bigcup_{i \in \mathbb{N}} B_i$$

is called as the *set generated by B*. Moreover if $B^* = A$, the B is called the *generator of A*.



Generators

- ▶ If \star is an operation such that inverses are well defined then in addition to elements of the form $b_i \star b_j$ we add b^{-1} as well in the process of “generation”.

Additive chains

An *additive chain* ending in a n is a sequence a_1, a_2, \dots, a_m such that for every $1 < i \leq m$, $a_i = a_j + a_k$ where $j, k < i$ and $a_m = n$. m is called the *length* of the chain

Open question

Given an n find the chain of smallest length ending in n .



Subgroup

- ▶ Let (A, \star) be a group and let $B \subseteq A$. B is a **subgroup** of A if B is a group.
- ▶ Given set B how do we check if it forms a subgroup of A ?
- ▶ We need to verify closure, associativity, existence identity and existence of inverse.
- ▶ Associativity comes for free.
- ▶ Identity of A must be the identity of B . (Why?)
- ▶ Inverse of an element in B must exist and be the same as its inverse in the group A . (Why?)



Subgroup

Subgroup criterion

A subset B of a group (A, \star) is a subgroup if and only if

- ▶ $B \neq \emptyset$
- ▶ For every $x, y \in B$, $x \star y^{-1} \in B$

Proof

- ▶ Since B is non empty there must exist an element x . Since $x \star x^{-1}$, the identity of A (therefore the identity B as well) must be present in B .
- ▶ Take x to be the identity of A and y any element of B . Thus $e \star y^{-1} = y^{-1}$ is present in B .
- ▶ Let $a, b \in B$. Take $x = a$ and $y = b^{-1}$. Since $(b^{-1})^{-1} = b$ we can conclude that $a \star (b^{-1})^{-1} = a \star b \in B$



Cyclic groups

- ▶ A group whose generator is a singleton set is called a **cyclic group**.

Lemma

Every finite cyclic group is commutative.

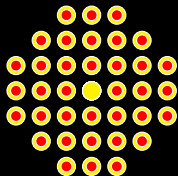
Proof

Let G be a finite cyclic group. Let g be the element in it's singleton generator. Therefore G is of the form $\{g^1, g^2, \dots, g^r\}$ where $g^r = \underbrace{g \star \dots \star g}_{r \text{ times}}$. The lemma follows from the fact that

$$g^i \star g^j = g^{i+j} = g^j \star g^i$$



An application of group theory.



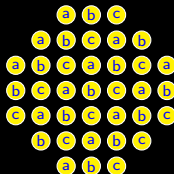
The diagram of left shows a solitaire game. The red circle inside a yellow circle denotes a position of the board with a marble.



- ▶ An allowed move is shown in the diagram above. A marble can jump over an adjacent marble (as indicated by the green arrow in figure).
- ▶ While jumping over a marble, one should remove that marble. Jumping may be done left to right, right to left, top to bottom and bottom to top.
- ▶ The resulting position after the jump in the example is shown on the right of the \Rightarrow .
- ▶ Can one reach a board configuration with a single marble?

Klein four group

- ▶ We shall show that it is impossible to reach a configuration with single coin.
- ▶ We will use the Klein four group (K_4) for this purpose.
- ▶ K_4 is defined on $\{a, b, c, e\}$. The operation \star is defined as
 - ▶ $a \star a = b \star b = c \star c = e$
 - ▶ K_4 is commutative.
 - ▶ e is the identity.
 - ▶ $a \star b = c, b \star c = a, c \star a = b$
- ▶ We will mark each yellow circle in the solitaire game described earlier by an element of K_4 as shown in figure above.
- ▶ Define **value** of a configuration to be product of elements at all the circles with marbles (red dots) in it.



Impossibility proof

- ▶ Note that the initial value of the board is $a^{12} \star b^{12} \star c^{12}$
- ▶ Each move changes the contents of exactly 3 yellow circles.
- ▶ Since any move involves 3 consecutive circles, they must be labeled using a , b and c .
- ▶ The contribution of these circles to the value of the board before the jump move is one among $\{a \star b = c, b \star c = a, c \star a = b\}$.
- ▶ Note that the contribution of these circles to the value of the board after the jump move is the same as its contribution before the jump move i.e. **Value of the board is invariant under allowable moves. It will be e .**
- ▶ No configuration having a single marble on board can have a board value of e . \square



Order

- ▶ The **order** of a group G denoted by $|G|$ or $\text{ord}\{G\}$ is the number of elements in the underlying set.
- ▶ The **order** of a element g of a group G denoted by $\text{ord}\{g\}$ is the smallest positive number n such that g^n is equal to the identity of G . When is is no such positive integer then we say that the order is infinite.
- ▶ The set $\{0, 1, \dots, 9\}$ forms a group under modulo 10 addition. The order of this group is 10. The order of 5 is 1 and the order of 3 is 10.



Lagrange's Theorem

Theorem

Let G be a finite group and H be a subgroup of G . $|H|$ divides $|G|$.

Proof

- ▶ Let $g \in G$. The subsets gH and Hg defined as below are the left and right coset of H w.r.t. g .

$$gH \triangleq \{gh \mid h \in H\}, \quad Hg \triangleq \{hg \mid h \in H\}$$

- ▶ Every coset of H has size $|H|$. (if $h_1 \neq h_2$ then $gh_1 \neq gh_2$.)
- ▶ If $g_1H \cap g_2H \neq \emptyset$, $\exists h_1, h_2 \in H$ such that $g_1h_1 = g_2h_2$.
- ▶ Since H is a group and $h_1, h_2 \in H$, $g_1 = g_2h_2h_1^{-1}$.
- ▶ $\therefore \forall h \in H, g_1h = g_2h_2h_1^{-1}h = g_2h_3$, for some $h_3 \in H$
- ▶ $\therefore g_1H \subseteq g_2H$. But as all cosets are of size $|H|$, $g_1H = g_2H$
- ▶ In other words, there cannot be overlapping cosets unless they are one and the same.



Lagrange's Theorem

Proof (Contd.)

- ▶ The union of all distinct cosets of H (overlapping cosets accounted only once) is G as H contains the identity.
- ▶ \therefore number of distinct cosets \times size of a coset = $|G|$
- ▶ Size of a coset = $|H|$. Thus we have

$$|H| = \frac{|G|}{\text{number of distinct cosets}}$$



Symmetric Group S_Ω

- ▶ Let Ω be any non empty set. Let S_Ω be the set of all one to one and onto functions (bijections) from Ω to itself. S_Ω forms a group under function composition.
- ▶ Suppose $\Omega = \{1, 2, \dots, n\} = [n]$. Then S_Ω is also referred to as the symmetric group of degree n written as S_n .
- ▶ A particular element of S_n can be written as below ($n = 8$).

$$\begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 6 & 3 & 2 & 4 & 5 & 1 & 8 & 7 \end{pmatrix}$$

- ▶ The above representation denotes a function σ where $\sigma(1) = 6, \sigma(2) = 3, \sigma(3) = 2, \sigma(4) = 4, \sigma(5) = 5, \sigma(6) = 1, \sigma(7) = 8$ and $\sigma(8) = 7$.



Example S_3

The elements of S_3 are as follows

$$\sigma_1 = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 3 \end{pmatrix} \sigma_2 = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 3 & 2 \end{pmatrix} \sigma_3 = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}$$

$$\sigma_4 = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 3 & 1 \end{pmatrix} \sigma_5 = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix} \sigma_6 = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 2 & 1 \end{pmatrix}$$

- ▶ Note that $\sigma_2(\sigma_3(1)) = 3, \sigma_2(\sigma_3(2)) = 1$ & $\sigma_2(\sigma_3(3)) = 2$
- ▶ If we denote the group operation by \circ , $\sigma_2 \circ \sigma_3 = \sigma_5$
- ▶ As $\sigma_3 \circ \sigma_2 = \sigma_4$, We know that S_3 is not an Abelian group.



Cycle Representation

- ▶ Instead of writing an element of S_n in two rows, we may represent it using cycles.
- ▶ Consider the permutation σ given below

$$\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 6 & 3 & 5 & 4 & 2 & 1 & 8 & 7 & 9 \end{pmatrix}$$

- ▶ $\sigma(1) = 6, \sigma(6) = 1$. (This completes a cycle)
- ▶ Further, $\sigma(2) = 3, \sigma(3) = 5, \sigma(5) = 2$
- ▶ Continuing this way, σ can be broken down into cycles and written as follows $(1, 6), (2, 3, 5), (4), (7, 8)(9)$
- ▶ Each (a_1, a_2, \dots, a_k) denotes a cycle such that $\sigma(a_i) = a_{i+1}$ for all i except k . and $\sigma(a_k) = a_1$.
- ▶ We will remove the cycles of length 1, for example in σ given above we shall remove the cycles (4) and (9)
- ▶ The representation obtained is called the cycle representation.



Cycle Representation (Algorithm)

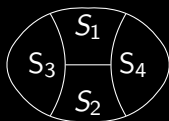
1. From $[n]$ pick the smallest element a which has not yet appeared in any cycle.
2. Add a as the starting element of a new cycle.
3. Compute the value $\sigma(x)$ where x the most recently added element of the cycle and repeat this till the cycle closes. Return to step 1 after cycle closes.
4. Remove cycles of length 1.

Remark: Let n be a permutation of $[n]$. $ord(n)$ will be equal to the l.c.m of the length of cycles in the cycles representation of σ .



Equivalence Relation and Partitions

- ▶ Given a set S , A **partition** of S is set of disjoint subsets of S such that their union is S .



- ▶ In Figure above, the set S is shown partitioned into 4 parts.
- ▶ An **equivalence relation** R is a binary relation defined on a set S such that S is
 - ▶ *reflexive*: aRa for all $a \in S$.
 - ▶ *symmetric*: If aRb then bRa .
 - ▶ *transitive*: If aRb and bRc then aRc .
- ▶ Partitions and Equivalence relations are one and the same thing.



Equivalence Relations and Partitions

- ▶ Given a partition P of S define a relation R such that aRb if and only if a and b belong to the same disjoint subset in the partition P . Verify that this relation R is indeed an equivalence relation.
- ▶ Conversely, given an equivalence relation R on a set S , we can define a partition in the following way.
 - ▶ For each element a define R_a to be the set of all elements b such that aRb . R_a is called the **equivalence class** of a .
 - ▶ Note that every element is in some equivalence class.
 - ▶ Also if two equivalence classes R_a and R_b have an overlap, then one can easily show that $R_a = R_b$ as R is an equivalence relation.
 - ▶ The set of all equivalence classes of elements in S thus forms a partition of S .



Homomorphism and Isomorphism

- ▶ Let (G, \star) and (H, \circ) be groups. Any function f from G to H such that for all $g_1, g_2 \in G$, $f(g_1 \star g_2) = f(g_1) \circ f(g_2)$ is called a **homomorphism**.
- ▶ For each $h \in H$, all the elements of G mapping to h forms the **fiber of f over h** .
- ▶ For example, the map $x \mapsto e^x$ is a homomorphism from $(\mathbb{R}, +)$ to (\mathbb{R}^+, \times) because $e^{x+y} = e^x e^y$. (\mathbb{R}^+ denotes the set of positive reals.)
- ▶ If f is a bijection then the homomorphism becomes an **isomorphism**.
- ▶ f maps the identity of G to the identity of H .
($f(e) \circ f(e) = f(e \star e) = f(e)$. Cancel $f(e)$ from both sides.)
- ▶ Image of the inverse equals the inverse of the image, i.e.
 $f(x^{-1}) = f(x)^{-1}$. (As $f(x) \circ f(x^{-1}) = f(x \star x^{-1}) = f(e)$)
- ▶ **Kernel** of f is the set of all a such that $f(a) = \text{identity of } H$.



More on Homomorphisms

Theorem

Kernel of f denoted by $\text{Ker}(f)$ is a subgroup of G and image of G under f is a subgroup of H .

Proof

For any $x, y \in G$ and a homomorphism f , we have $f(xy^{-1}) = f(x)f(y^{-1}) = f(x)f(y)^{-1}$. Let e_1 be the identity in G and e_2 the identity in H

Consider $x, y \in \text{Ker}$. Thus $f(x) = f(y) = f(y)^{-1} = e_1$.

$\therefore f(xy^{-1}) = f(x)f(y)^{-1} = e$. Thus $\text{Ker}(f)$ is a group.

Consider $x, y \in \text{Image of } G$ Thus $\exists x', y' \text{ in } G$ such that $f(x') = x$ and $f(y') = y$. Note that $f(x'y'^{-1}) = xy^{-1}$ Thus $xy^{-1} \in \text{Image of } G$. Thus image of G is a group.

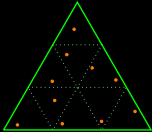


Quotient Group

- ▶ Let f be a homomorphism with K as its kernel. The **quotient group** G/K (read as $G \bmod K$) is the group consisting of fibers of f .
- ▶ Suppose G_1 is the fiber above $f(g_1)$ and G_2 the fiber above $f(g_2)$, then the product in G/K is computed by defining $G_1 G_2$ to be the fiber above $f(g_1 g_2)$. (Verify that this definition is well defined as f is homomorphism)
- ▶ $GL_n(\mathbb{R})$ be the group of all invertible $n \times n$ matrices. Consider the map from $GL_n(\mathbb{R})$ to $\mathbb{R} \setminus \{0\}$ given by $f : A \mapsto \det(A)$ where $\det(A)$ denotes the determinant of A .
- ▶ $\text{Ker}(f)$ is all $n \times n$ matrices with determinant one. Fibers are all matrices with determinant $c, c \in \mathbb{R} \setminus \{0\}$



Pigeon Hole Principle



Let 10 points be randomly chosen from an equilateral triangle of side 3. Show that there will be two points within distance 1 cm of each other.

- ▶ Observe that there must be at least one small triangle which contains 2 points.
- ▶ These points must be within 1 cm of each other.

Pigeon Hole Principle

Suppose there are n objects to be distributed into $n - 1$ boxes, then there exists a box which contains more than one object.

Generalized Pigeon Hole Principle

Suppose there are $q_1 + q_2 + \dots + q_n - n + 1$ objects to be distributed into n boxes, then there exists an i such that the i th box contains at least q_i objects in it.

Applications of PHP

- ▶ In every set of people (with more than two people) there exists two persons with same number of friends.

Proof

- ▶ Let the set have n people.
- ▶ If there is a person with no friends there cannot be a person who friends with everyone.
- ▶ If there is a person who is friends with everyone there cannot be a person without friends.
- ▶ Thus the number of friends each person can have is either from the set $\{0, 1, \dots, n - 2\}$ or from the set $\{1, \dots, n - 1\}$. In either case the number of distinct elements in the set is $n - 1$. Therefore one element must repeat. \square



Applications of PHP

Erdős Szekeres Theorem

Every sequence of length $n^2 + 1$ contains a monotone subsequence of length $n + 1$.

Proof

- ▶ Let $a_1, a_2, \dots, a_{n^2+1}$ be the series we are considering.
- ▶ Let m_i denote the length of the maximal increasing sequence starting from the element a_i
- ▶ If any m_i is greater than n we have a monotone subsequence of length $n + 1$. So let us assume that every m_i is less than or equal to n .
- ▶ Since there are $n^2 + 1$ different m_i , taking values 1 to n , there must exist an L such that $n + 1$ of the m_i s takes the value L .



Erdős Szekeres Theorem (Proof)

- ▶ Consider the a_i 's with $m_i = L$.
- ▶ Let us write these elements as b_1, b_2, \dots, b_k (Note that $k \geq n + 1$). without changing the order in which they appear in the original sequence.
- ▶ If $b_s < b_{s+1}$ for any s , then consider the maximal monotonic increasing sequence starting at b_{s+1} (which is of length L). We can append b_s to the start of this sequence to get an increasing monotonic sequence of length $L + 1$ starting at b_s .
- ▶ This contradicts the assumption that the maximal increasing monotonic sequence starting at b_s is of length L .
- ▶ Thus we have $b_s \geq b_{s+1}$ for all s .
- ▶ Considering b_i 's, we have obtained a monotonic decreasing sequence of length at least $n + 1$ \square



Principle of Mathematical Induction

Weak Induction

Let $P(n)$ be statement about a natural number n such that

- ▶ **Base:** $P(1)$ is true.
- ▶ **Induction:** $P(n + 1)$ true whenever $P(n)$ is true

Then $P(n)$ is true for all $n \in \mathbb{N}$

Strong Induction

Let $P(n)$ be statement about a natural number n such that

- ▶ **Base:** $P(1)$ is true.
- ▶ **Induction:** $P(n + 1)$ true whenever $P(m)$ is true for $m \leq n$

Then $P(n)$ is true for all $n \in \mathbb{N}$



Incorrect use of PMI

(Pseudo)Theorem: All horses are of the same color.

(Pseudo) Proof:

$P(n) \triangleq$ Any set of containing n horses have horses of identical color.

- ▶ Base case: For $n = 1$ the statement $P(n)$ is certainly true.
- ▶ Induction case: Assume that $P(k)$ is true for some k . Now consider a set of $k + 1$ horses.
- ▶ The horses numbered 1 to k forms a set of k horses. They are all of the same color say c . In particular the horse numbered k is of color c .
- ▶ The horses numbered 2 to $k + 1$ forms a set of k horses. They are all of the same color as the horse numbered k i.e. c . Thus all the horses are of the same color.

Question: Where is the mistake in the above “proof”?



Well Ordering Principle

Every non empty subset of \mathbb{N} has a smallest element.

- ▶ Well ordering principle is equivalent to PMI.
- ▶ We shall first prove that PMI \Rightarrow WOP using strong induction.

$P(n) \triangleq$ Every subset of \mathbb{N} containing n has a least element

- ▶ **Base:** 1 is certainly the least element of any subset of \mathbb{N} containing 1. Thus $P(1)$ is true.
- ▶ **Induction:** Consider any set S containing $k + 1$.
- ▶ If S contains any element, say m , smaller than $k + 1$, then by strong induction, as $P(m)$ is true, we know that S contains a least element.
- ▶ If S didn't contain any element smaller than $k + 1$, then S contains a smallest element, namely $k + 1$. Thus $P(k + 1)$ is true. \square



Well Ordering Principle

- ▶ We shall now show the reverse direction namely WOP \Rightarrow PMI.
- ▶ For contradiction, let us assume that there is a property P such that
 - ▶ $P(1)$ is true and whenever $P(k)$ is true, $P(k + 1)$ is also true.
 - ▶ There exists a number m such that $P(m)$ is false.
- ▶ Let $S \triangleq \{x \in \mathbb{N} \mid P(x) \text{ is false}\}$.
- ▶ Since $m \in S$, S is a non empty subset of \mathbb{N} and thus has a least element say s .
- ▶ $s \neq 1$ because $P(1)$ is true. Since s is the least element of S , $s - 1 \notin S$.
- ▶ $\therefore P(s - 1)$ is true. But then $P((s - 1) + 1)$ must also be true and thus $s \notin S$. \square



Topics

- ▶ Introduction Definitions
- ▶ Eulerian Cycles
- ▶ Hamiltonian Cycles
- ▶ Tournament Graphs
- ▶ Minimal Spanning Trees

